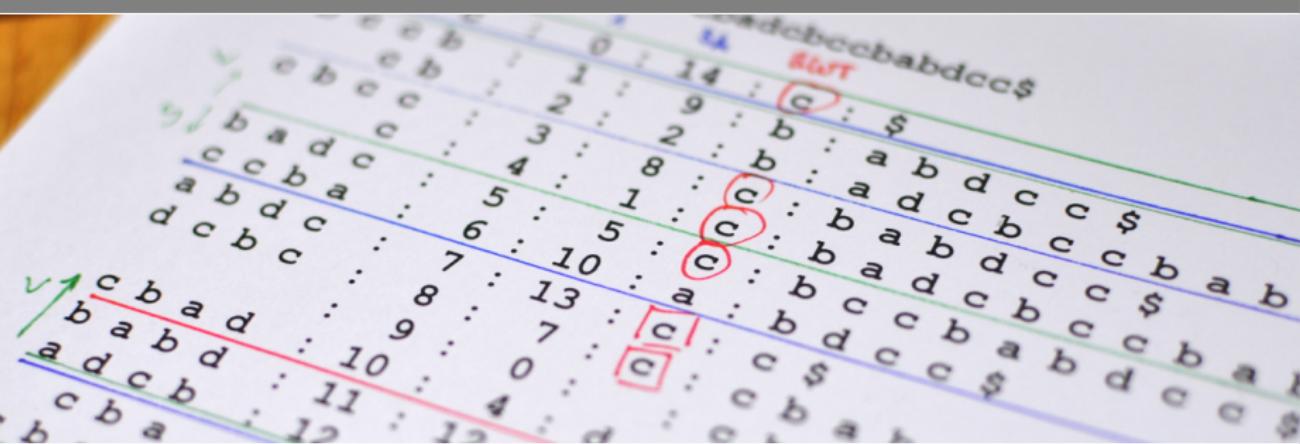


Inducing Suffix and LCP Arrays in External Memory

Timo Bingmann, Johannes Fischer, and Vitaly Osipov | September 5th, 2013 @ MASSIVE'13

INSTITUTE OF THEORETICAL INFORMATICS – ALGORITHMIC



Example $T = [\text{cababcbababb\$}]$

i	T_i
0	c a b a b c b a b a b b \\$
1	a b a b c b a b a b b \\$
2	b a b c b a b a b b \\$
3	a b c b a b a b b \\$
4	b c b a b a b b \\$
5	c b a b a b b \\$
6	b a b a b b \\$
7	a b a b b \\$
8	b a b b \\$
9	a b b \\$
10	b b \\$
11	b \\$
12	\\$

Example $T = [\text{cababcbababb\$}]$

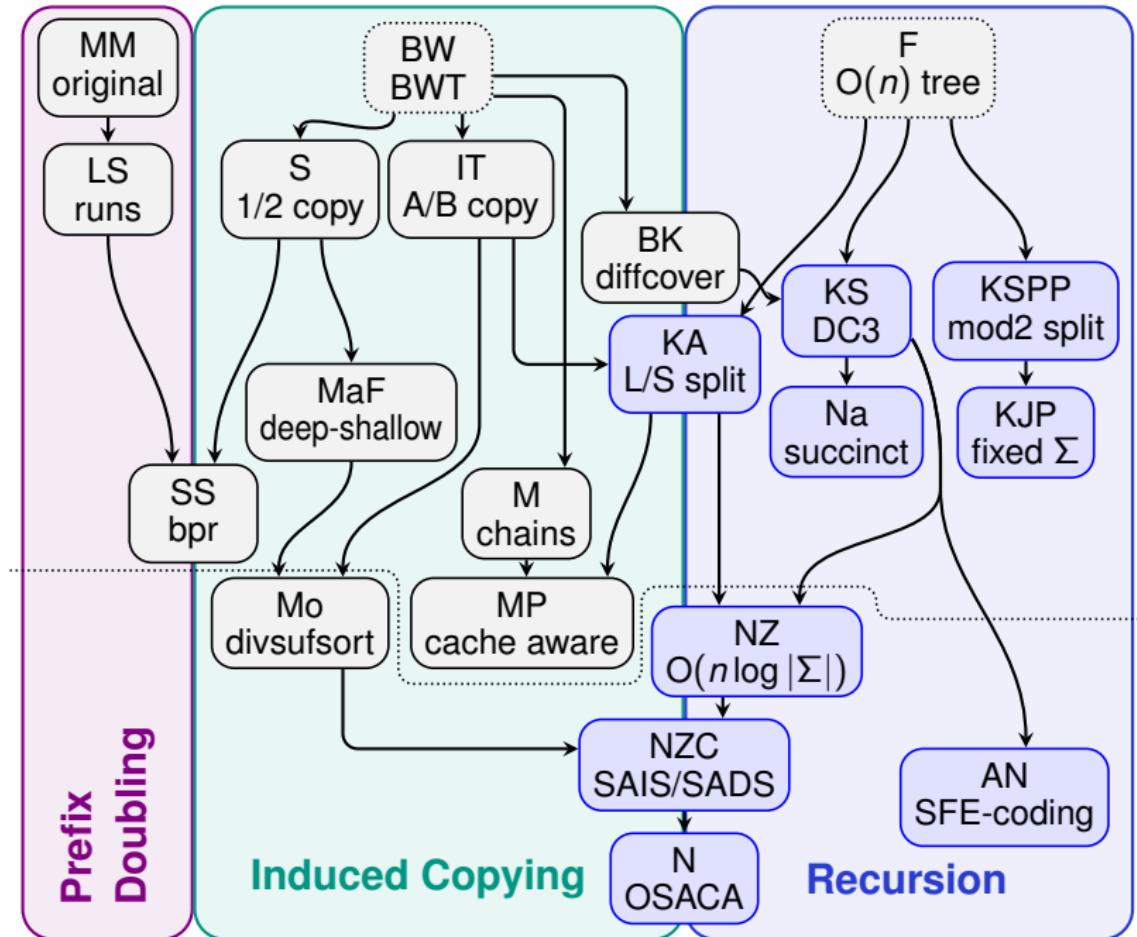
SA_i	$T_{SA_i \dots n}$
12	\$
7	a b a b b \$
1	a b a b c b a b a b b \$
9	a b b \$
3	a b c b a b a b b \$
11	b \$
6	b a b a b b \$
8	b a b b \$
2	b a b c b a b a b b \$
10	b b \$
4	b c b a b a b b \$
0	c a b a b c b a b a b b \$
5	c b a b a b b \$

Example $T = [\text{cababcbababb\$}]$

SA_i	LCP_i	$T_{\text{SA}_i \dots n}$
12		\$
7		a b a b b \$
1		a b a b c b a b a b b \$
9		a b b \$
3		a b c b a b a b b \$
11		b \$
6		b a b a b b \$
8	3	b a b b \$
2		b a b c b a b a b b \$
10		b b \$
4		b c b a b a b b \$
0		c a b a b c b a b a b b \$
5		c b a b a b b \$

Example $T = [\text{cababcbababb\$}]$

SA_i	LCP_i	$T_{\text{SA}_i \dots n}$
12	-	\$
7	0	a b a b b \$
1	4	a b a b c b a b a b b \$
9	2	a b b \$
3	2	a b c b a b a b b \$
11	0	b \$
6	1	b a b a b b \$
8	3	b a b b \$
2	3	b a b c b a b a b b \$
10	1	b b \$
4	1	b c b a b a b b \$
0	0	c a b a b c b a b a b b \$
5	1	c b a b a b b \$



[PST07]
updated

1999
2000

2003

2004

2005

2006

2007

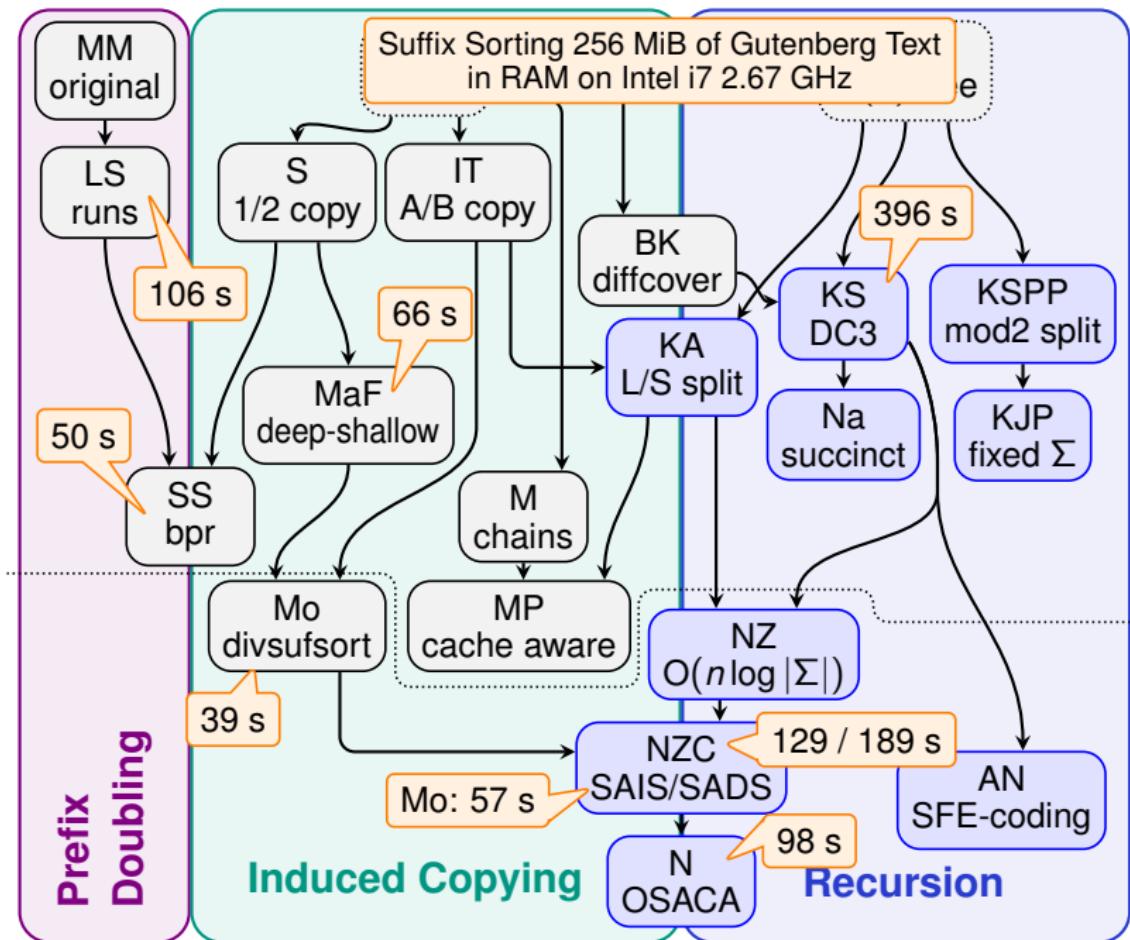
2009

2011

**Prefix
Doubling**

Induced Copying

Recursion



[PST07] updated

1999
2000

2003

2004

2005

2006

2007

2009

2011

[PST07]
updated

1999
2000

2003

2004

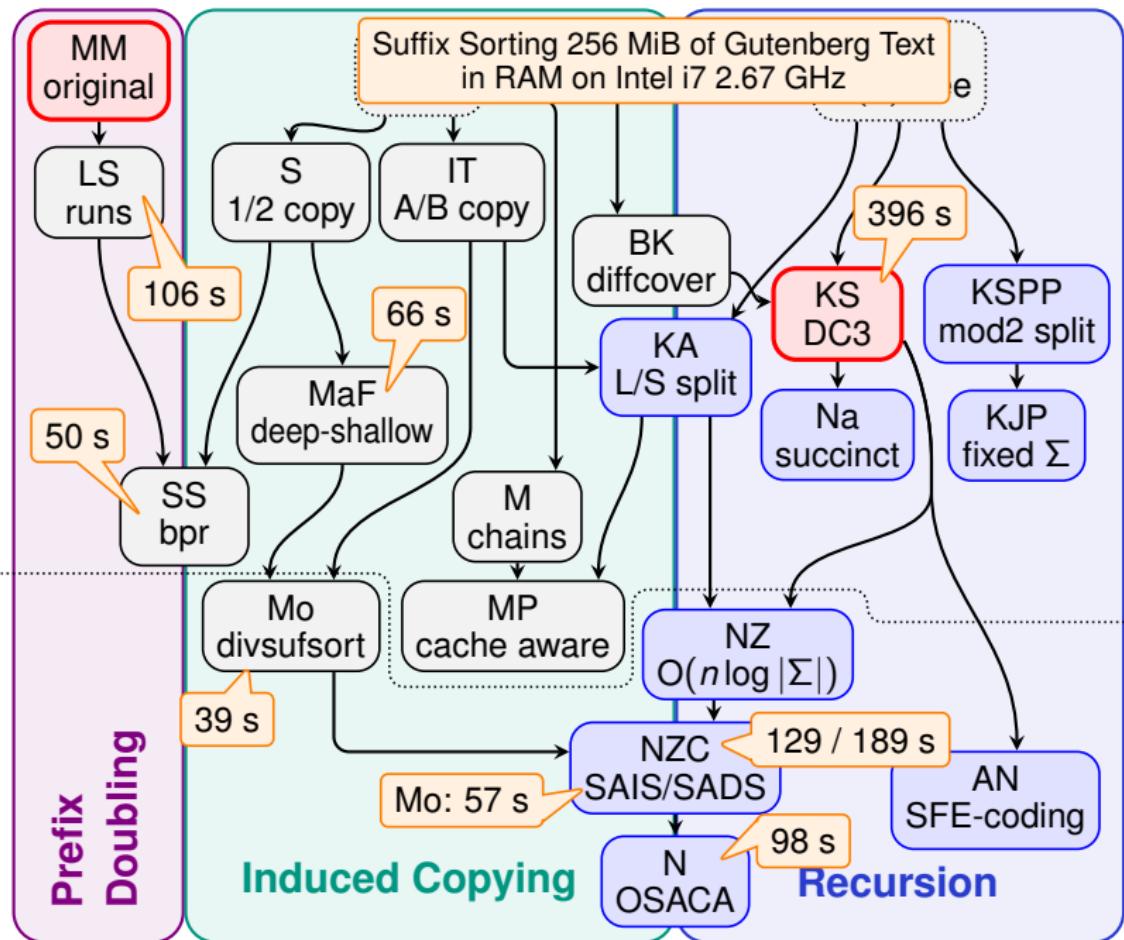
2005

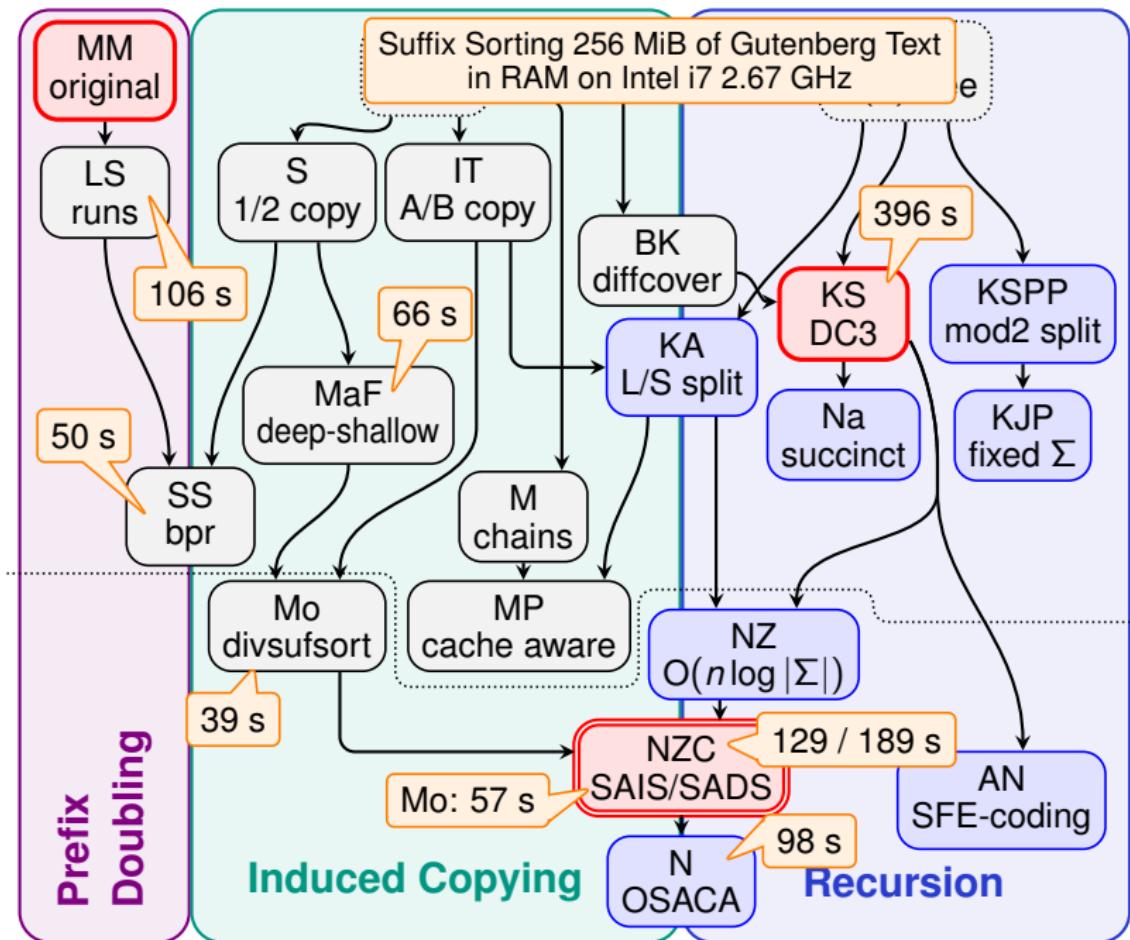
2006

2007

2009

2011





[PST07]
updated

1999
2000

2003

2004

2005

2006

2007

2009

2011

Example $T = [\text{cababcbababb\$}]$

SA_i	$T_{SA_i \dots n}$
12	\$
7	a b a b b \$
1	a b a b c b a b a b b \$
9	a b b \$
3	a b c b a b a b b \$
11	b \$
6	b a b a b b \$
8	b a b b \$
2	b a b c b a b a b b \$
10	b b \$
4	b c b a b a b b \$
0	c a b a b c b a b a b b \$
5	c b a b a b b \$

Example $T = [\text{cababcbababb\$}]$

SA_i	$T_{SA_i \dots n}$
12	\$
7	a b a b b \$
1	a b a b c b a b a b b \$
9	a b b \$
3	a b c b a b a b b \$
11	b \$
6	b a b a b b \$
8	b a b b \$
2	b a b c b a b a b b \$
10	b b \$
4	b c b a b a b b \$
0	c a b a b c b a b a b b \$
5	c b a b a b b \$

Example $T = [\text{cababcbababb\$}]$

SA_i	$T_{SA_i \dots n}$
12	\$
7	a b a b b \$
1	a b a b c b a b a b b \$
9	a b b \$
3	a b c b a b a b b \$
11	b \$
6	b a b a b b \$
8	b a b b \$
2	b a b c b a b a b b \$
10	b b \$
4	b c b a b a b b \$
0	c a b a b c b a b a b b \$
5	c b a b a b b \$

Example $T = [\text{cababcbababb\$}]$

SA_i	T_{i-1}	$T_{\text{SA}_i \dots n}$
12	b	\$
7	b	a b a b b \$
1	c	a b a b c b a b a b b \$
9	b	a b b \$
3	b	a b c b a b a b b \$
11	b	\$
6	c	b a b a b b \$
8	a	b a b b \$
2	a	b a b c b a b a b b \$
10	a	b b \$
4	a	b c b a b a b b \$
0	-	c a b a b c b a b a b b \$
5	b	c b a b a b b \$

Example $T = [\text{cababcbababb\$}]$

SA_i	T_{i-1}	$T_{\text{SA}_i \dots n}$
12	b	\$
7	b	a b a b b \$
1	c	a b a b c b a b a b b \$
9	b	a b b \$
3	b	a b c b a b a b b \$
11	b	b \$
6	c	b a b a b b \$
8	a	b a b b \$
2	a	b a b c b a b a b b \$
10	a	b b \$
4	a	b c b a b a b b \$
0	-	c a b a b c b a b a b b \$
5	b	c b a b a b b \$

Example $T = [\text{cababcbababb\$}]$

SA_i	T_{i-1}	$T_{\text{SA}_i \dots n}$
12	b \$	
7	b a b a b b \$	
1	c a b a b c b a b a b b \$	
9	b a b b \$	
3	b a b c b a b a b b \$	
11	b b \$	
6	c b a b a b b \$	
8	a b a b b \$	
2	a b a b c b a b a b b \$	
10	a b b \$	
4	a b c b a b a b b \$	
0	- c a b a b c b a b a b b \$	
5	b c b a b a b b \$	

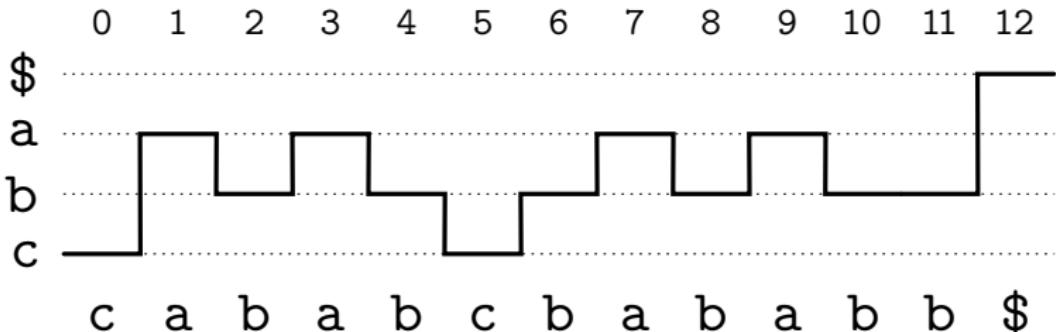
Example $T = [\text{cababcbababb\$}]$

SA_i	T_{i-1}	$T_{\text{SA}_i \dots n}$
12	b \$ ← 0	
7	b a b a b b \$ ← 1	
1	c a b a b c b a b a b b \$ ← 2	
9	b a b b \$ ← 3	
3	b a b c b a b a b b \$ ← 4	
11	b b \$	
6	c b a b a b b \$	
8	a b a b b \$	
2	a b a b c b a b a b b \$	
10	a b b \$	
4	a b c b a b a b b \$	
0	- c a b a b c b a b a b b \$	
5	b c b a b a b b \$	

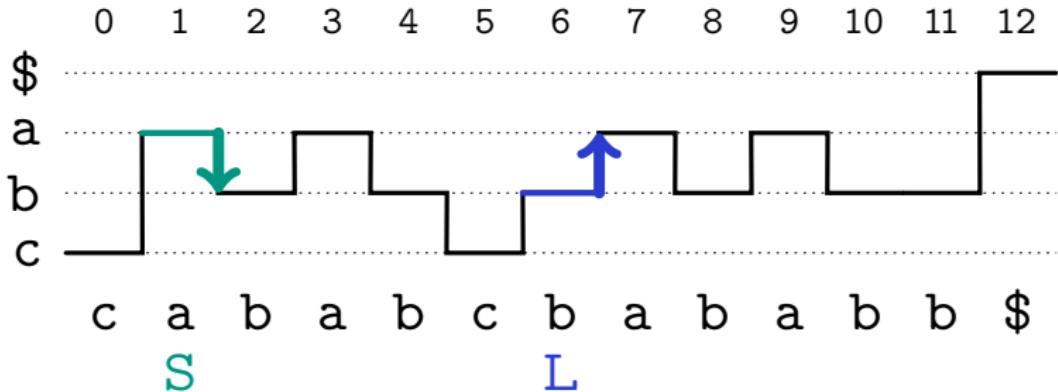
Example $T = [\text{cababcbababb\$}]$

SA_i	T_{i-1}	$T_{\text{SA}_i \dots n}$
12	b	\$ ← b, 0
7	b a b a b b \$	← b, 1
1	c a b a b c b a b a b b \$	← c, 2
9	b a b b \$	← b, 3
3	b a b c b a b a b b \$	← b, 4
11	b b \$	
6	c b a b a b b \$	
8	a b a b b \$	
2	a b a b c b a b a b b \$	
10	a b b \$	
4	a b c b a b a b b \$	
0	- c a b a b c b a b a b b \$	
5	b c b a b a b b \$	

Example $T = [\text{cababcbababb\$}]$

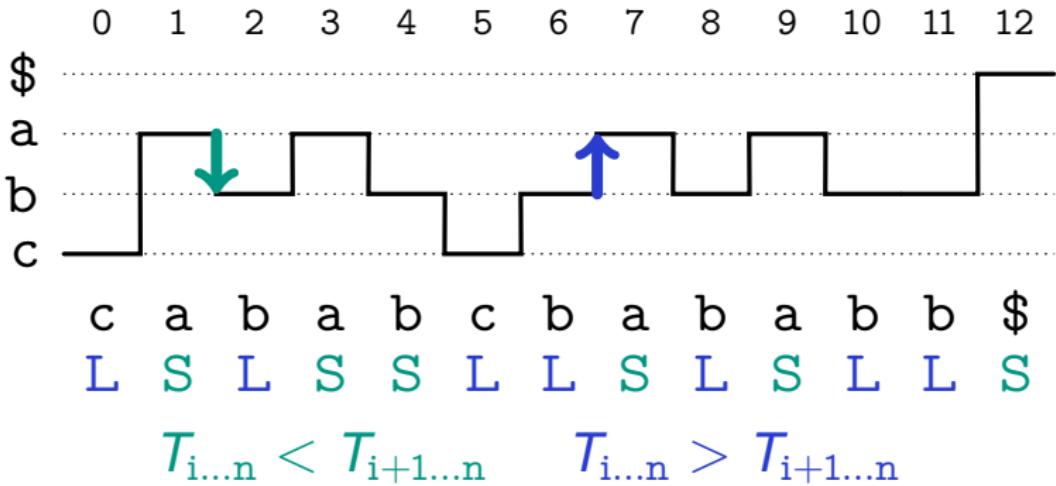


Example $T = [\text{cababcbababb\$}]$

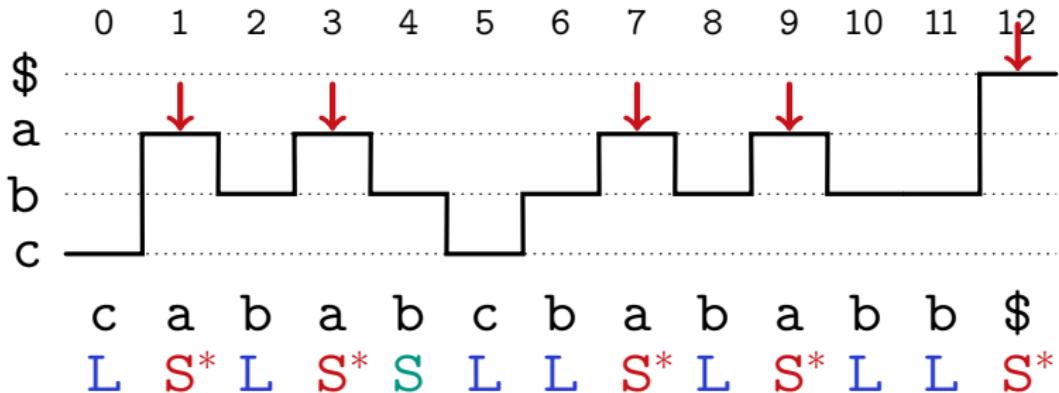


$$T_{i \dots n} < T_{i+1 \dots n} \quad T_{i \dots n} > T_{i+1 \dots n}$$

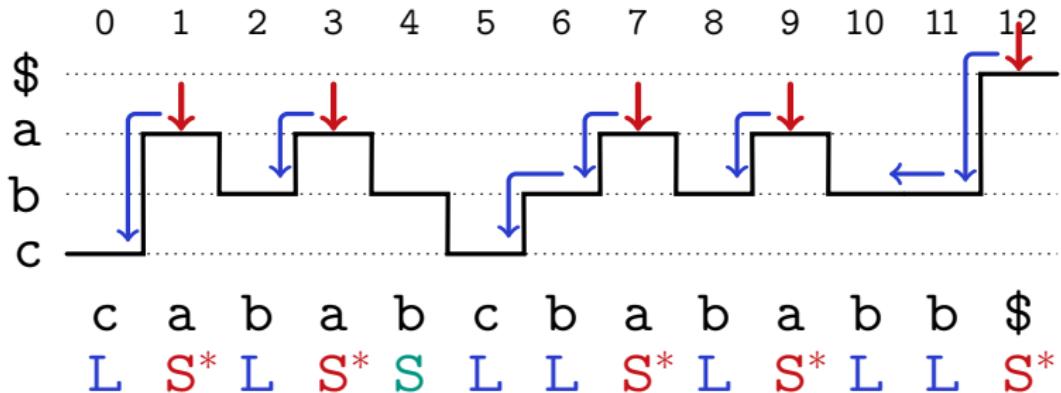
Example $T = [\text{cababcbababb\$}]$



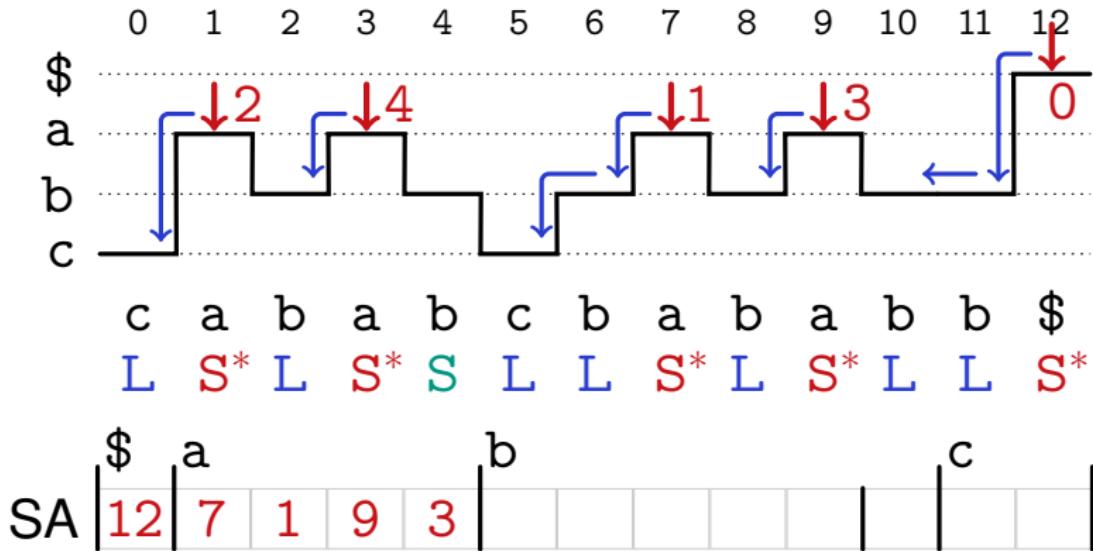
Example $T = [\text{cababcbababb\$}]$



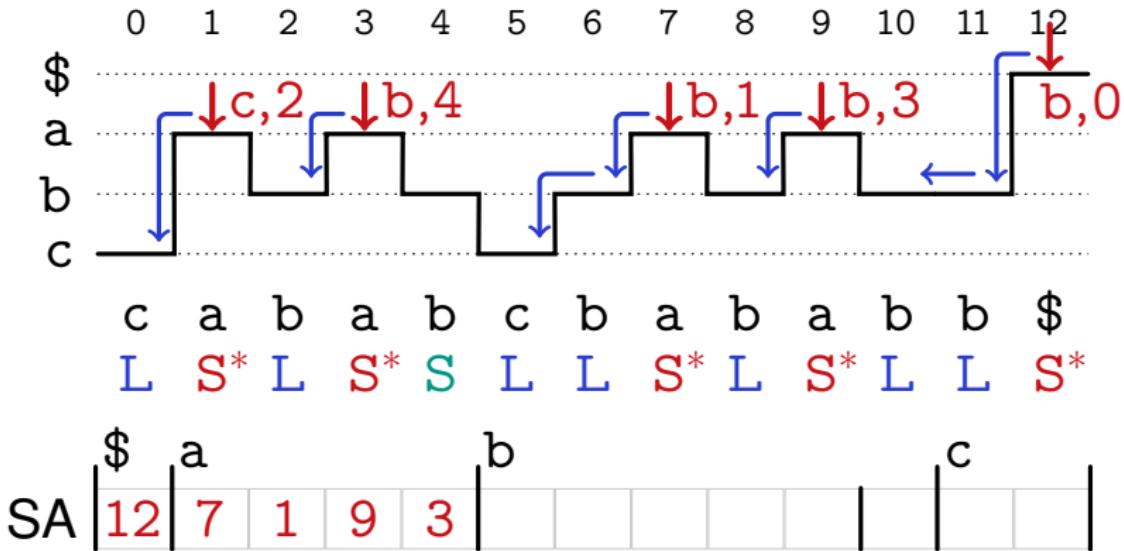
Example $T = [\text{cababcbababb\$}]$



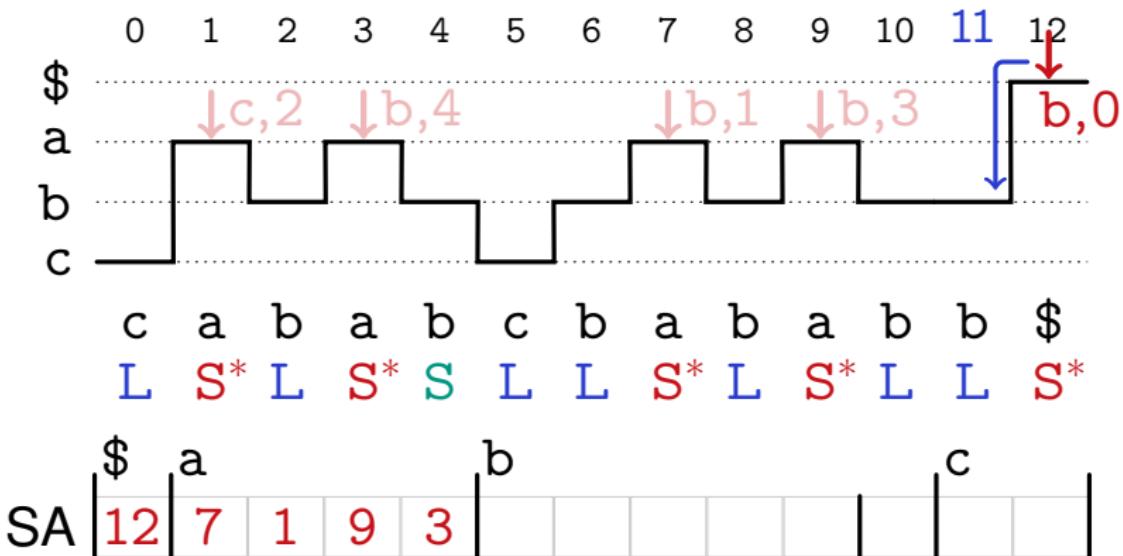
Example $T = [\text{cababcbababb\$}]$



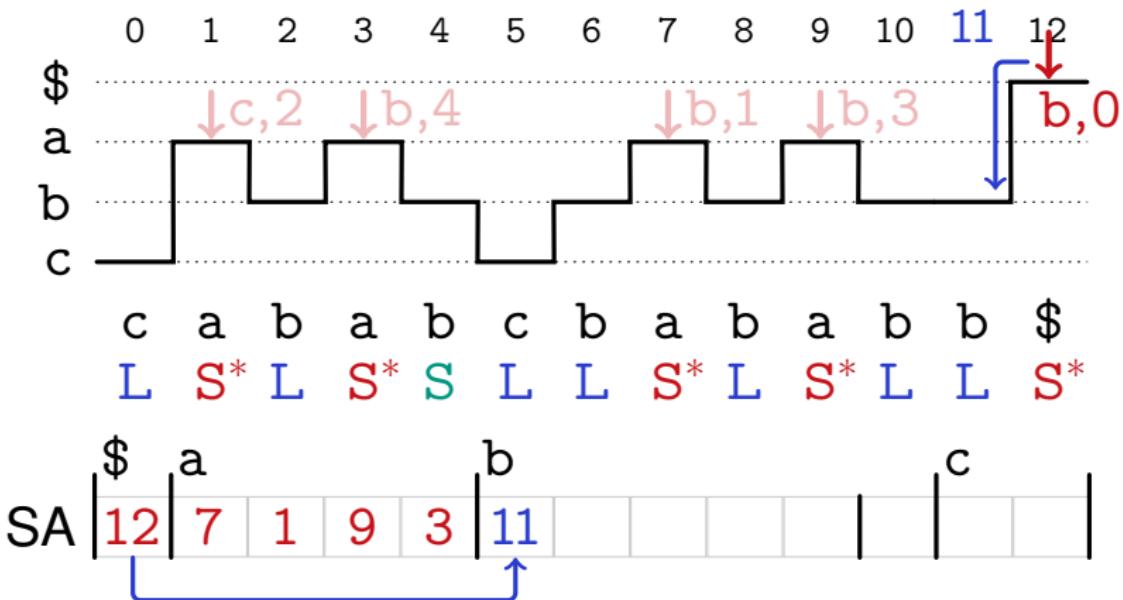
Example $T = [\text{cababcbababb\$}]$



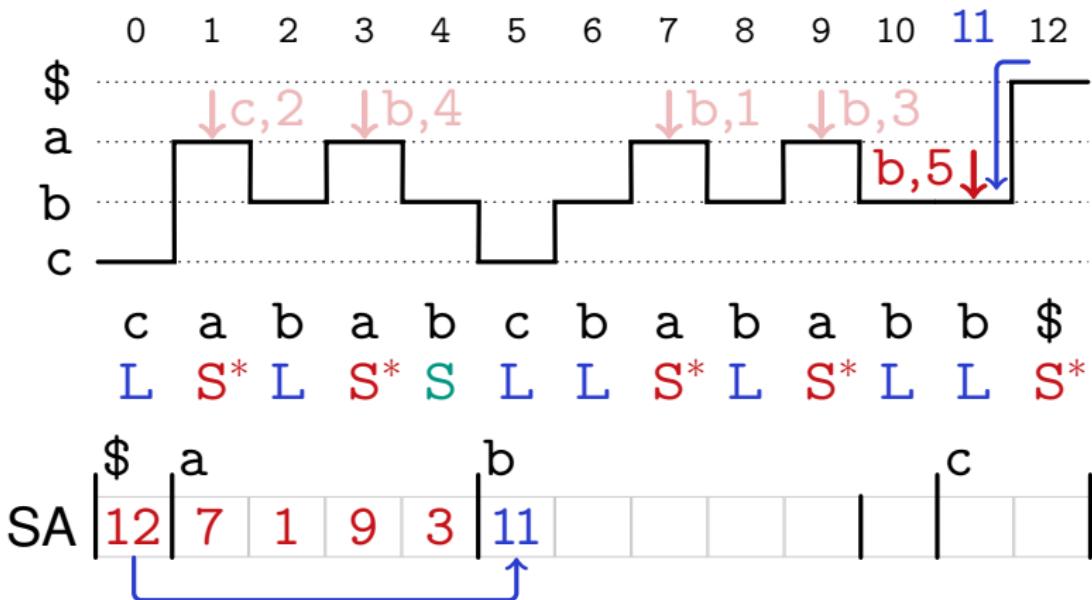
Example $T = [\text{cababcbababb\$}]$



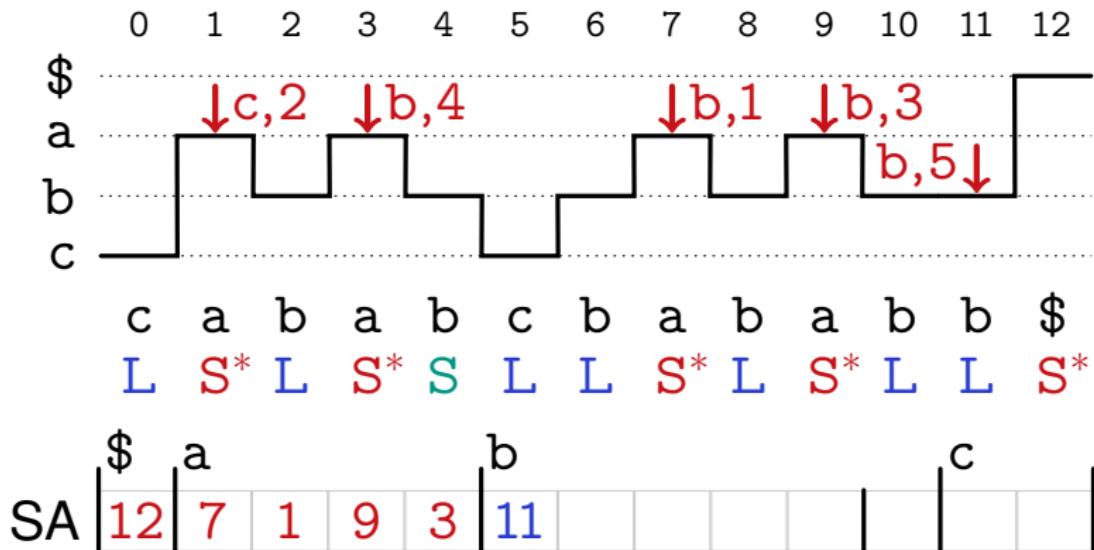
Example $T = [\text{cababcbababb\$}]$



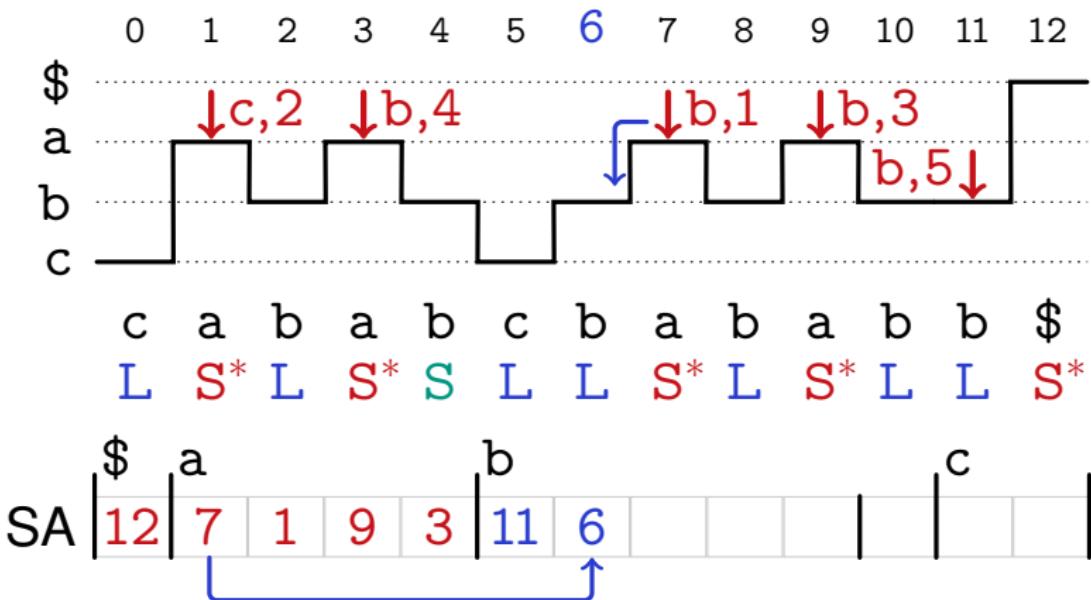
Example $T = [\text{cababcbababb\$}]$



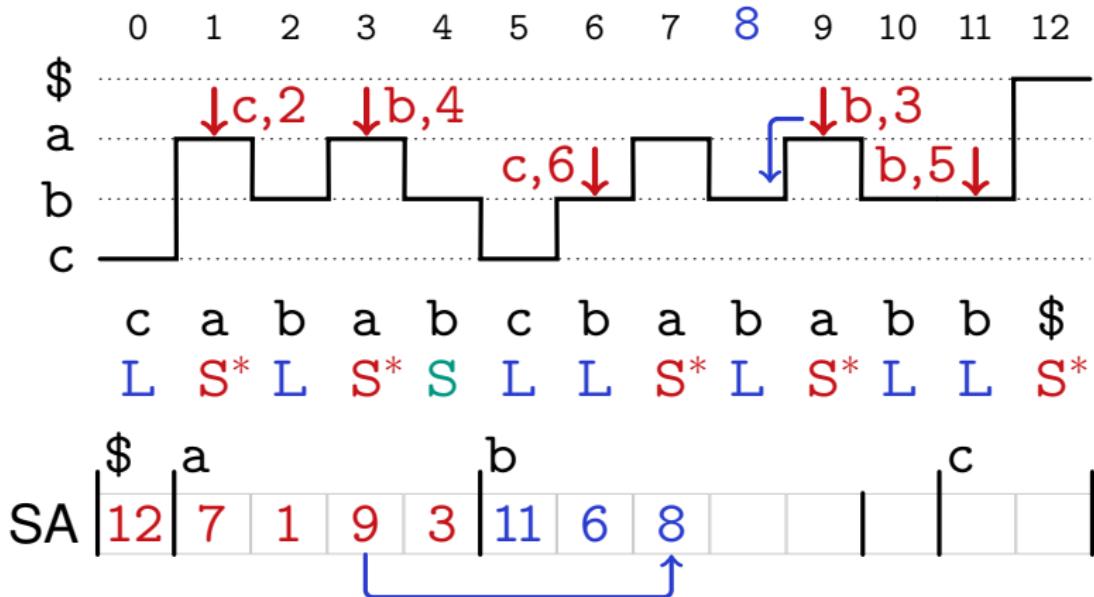
Example $T = [\text{cababcbababb\$}]$



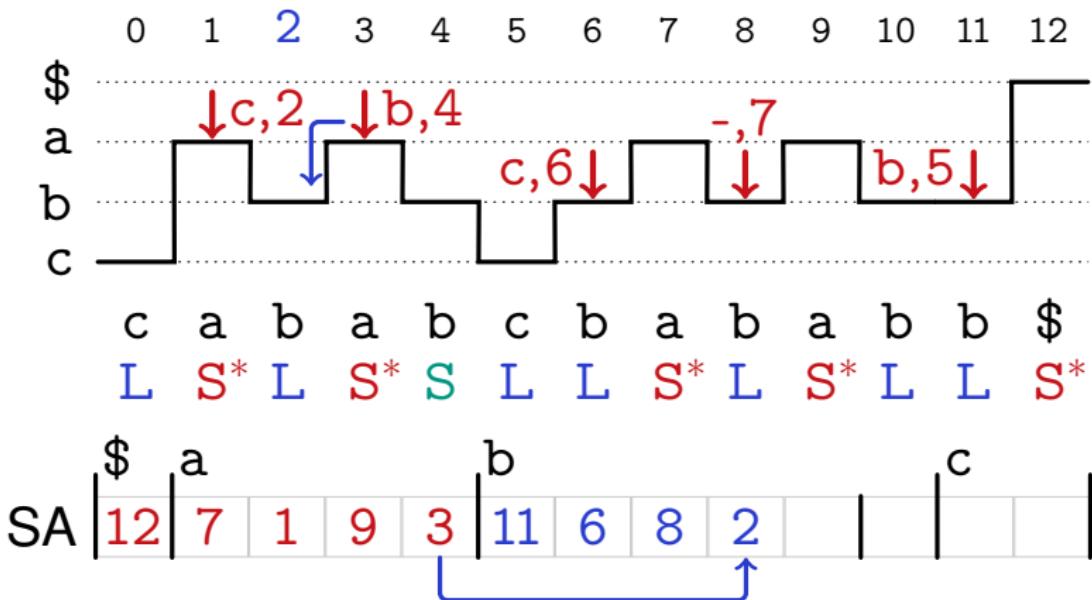
Example $T = [\text{cababcbababb\$}]$



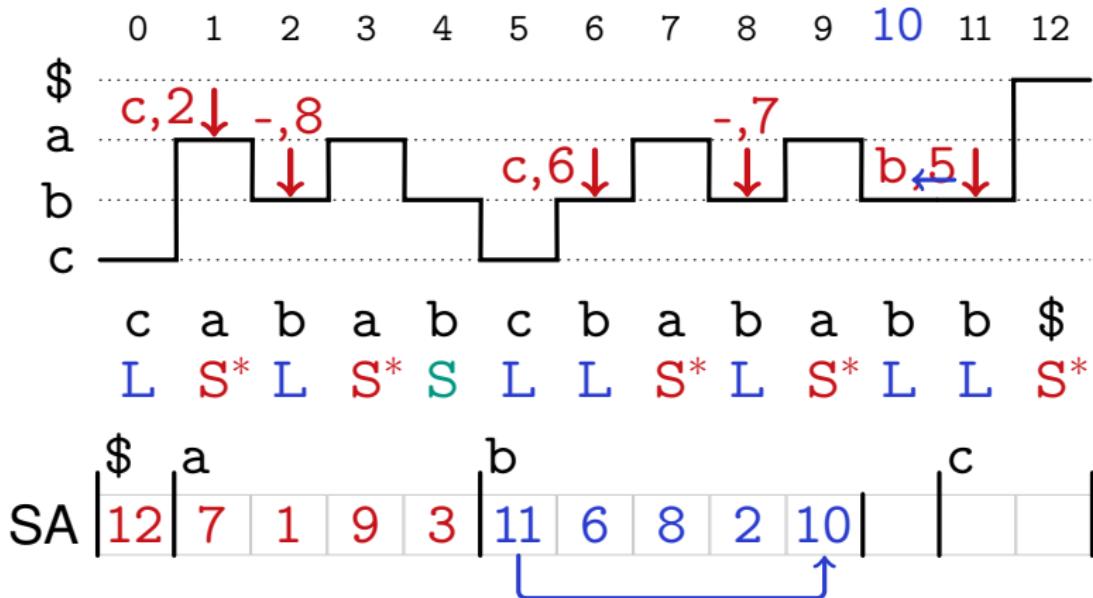
Example $T = [\text{cababcbababb\$}]$



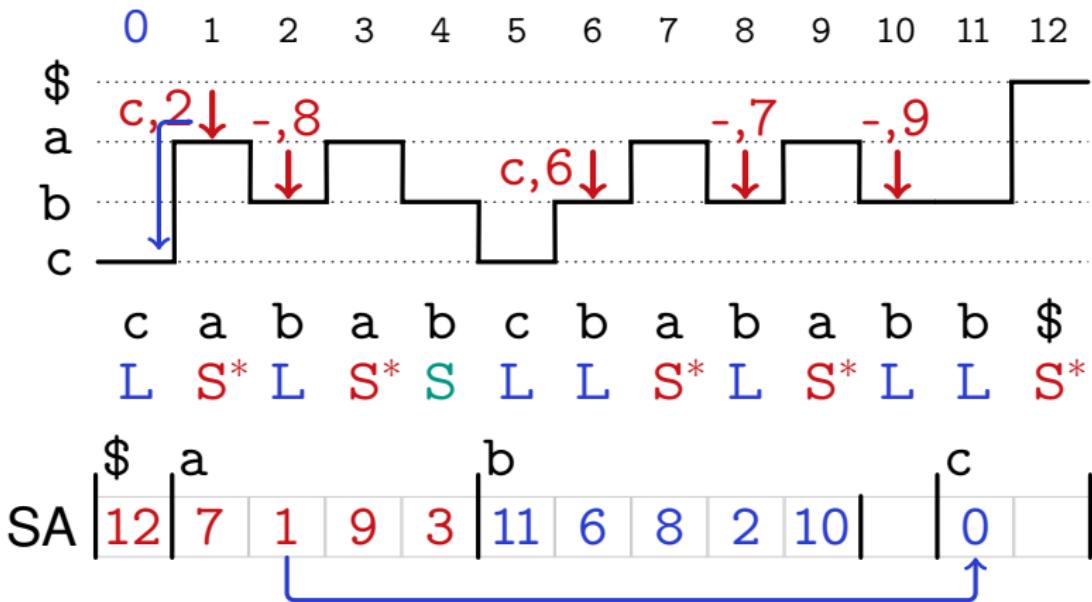
Example $T = [\text{cababcbababb\$}]$



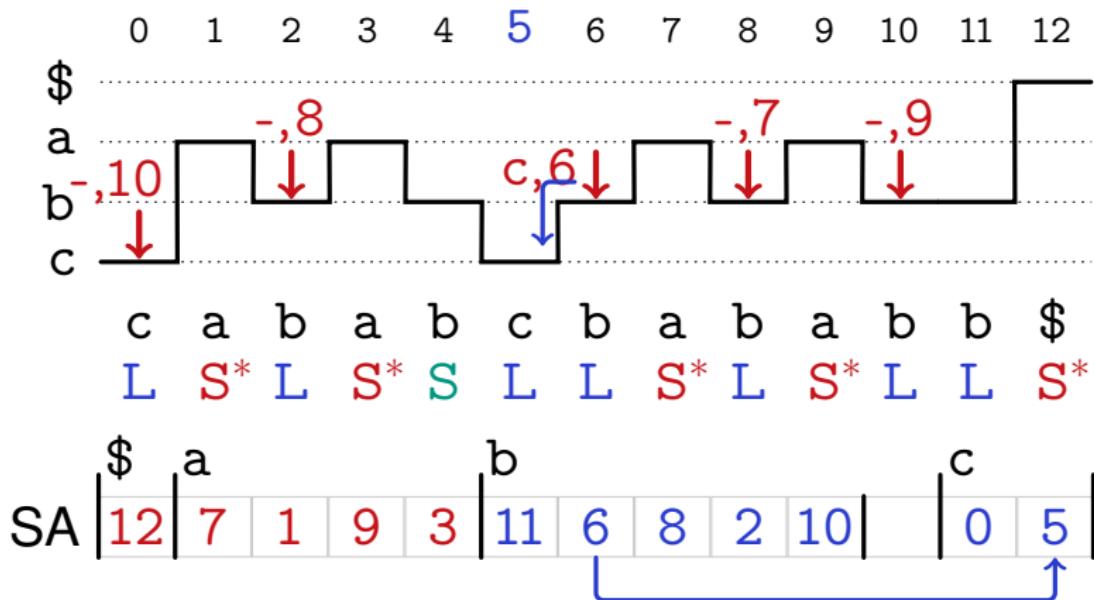
Example $T = [\text{cababcbababb\$}]$



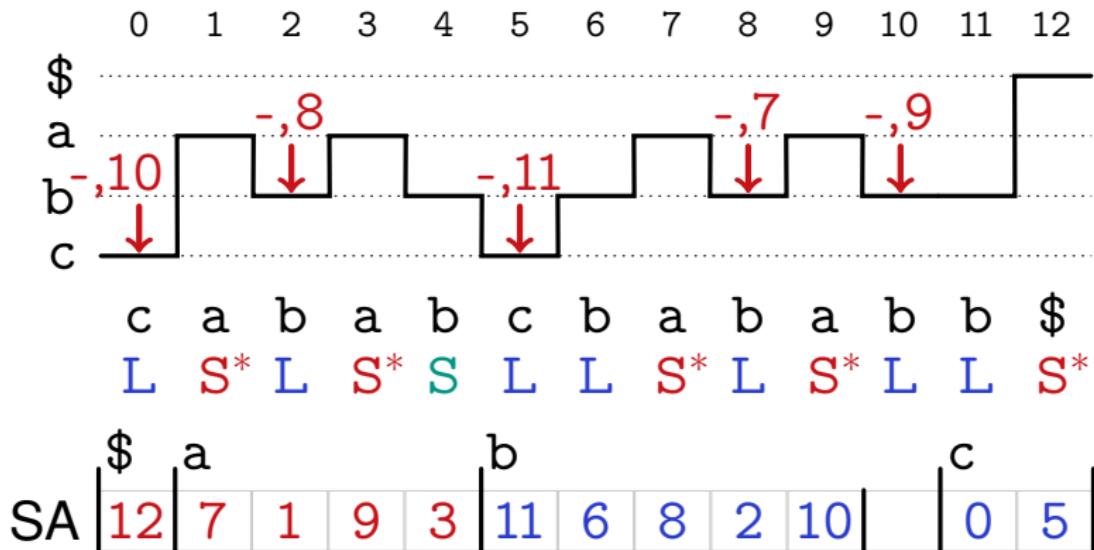
Example $T = [\text{cababcbababb\$}]$



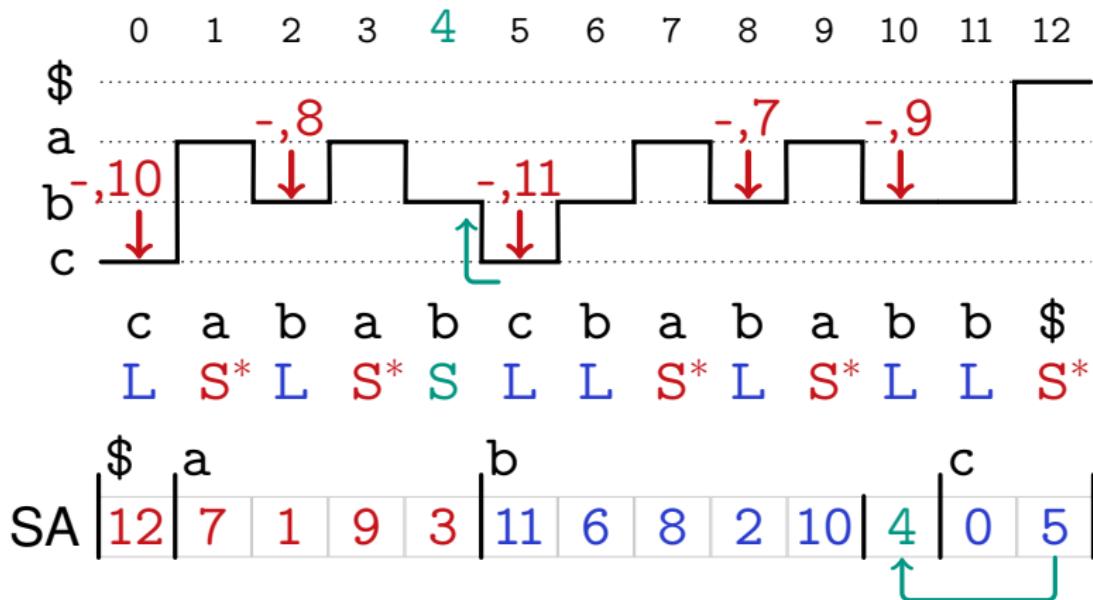
Example $T = [\text{cababcbababb\$}]$



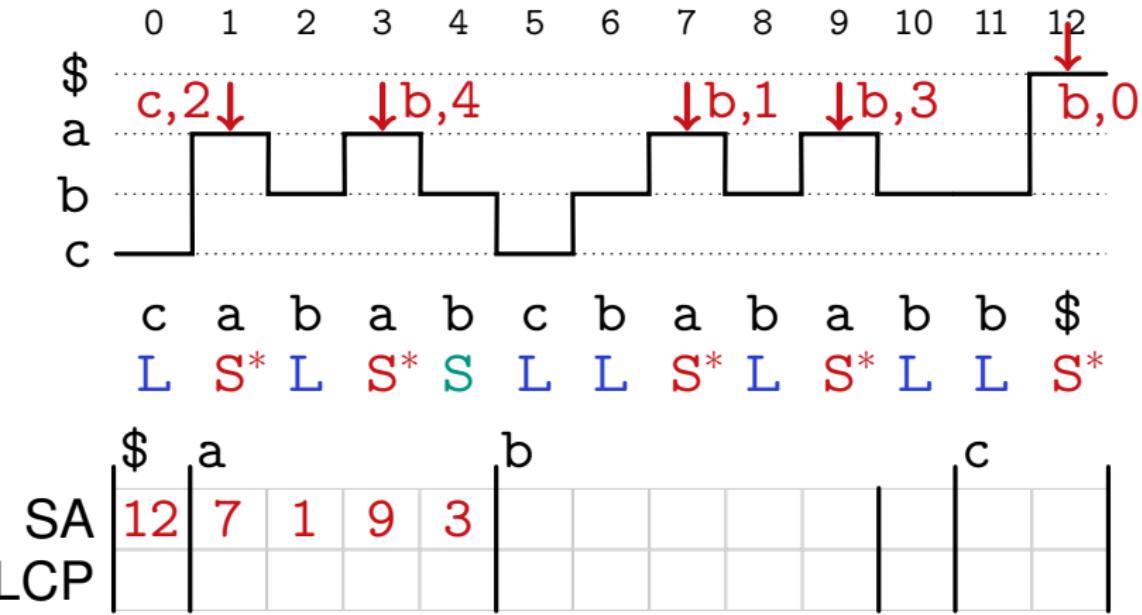
Example $T = [\text{cababcbababb\$}]$



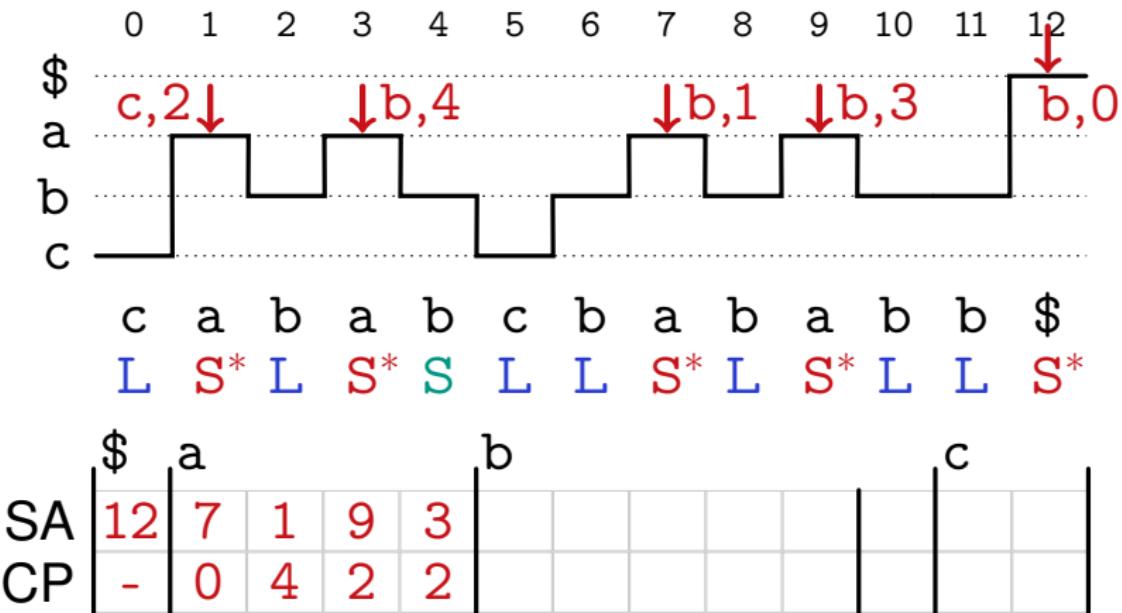
Example $T = [\text{cababcbababb\$}]$



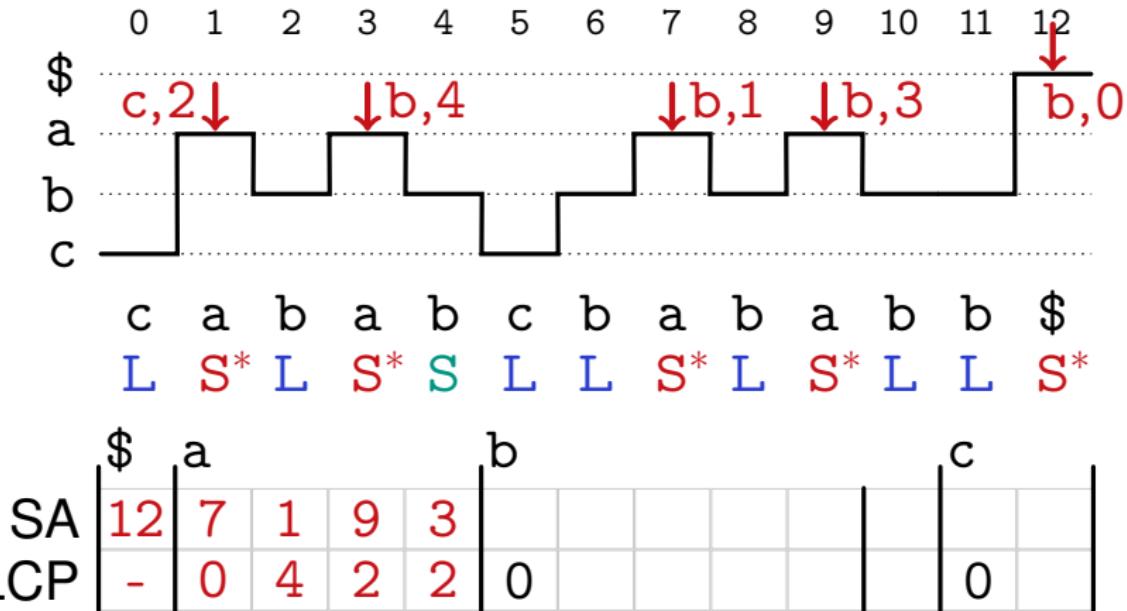
Example $T = [\text{cababcbababb\$}]$



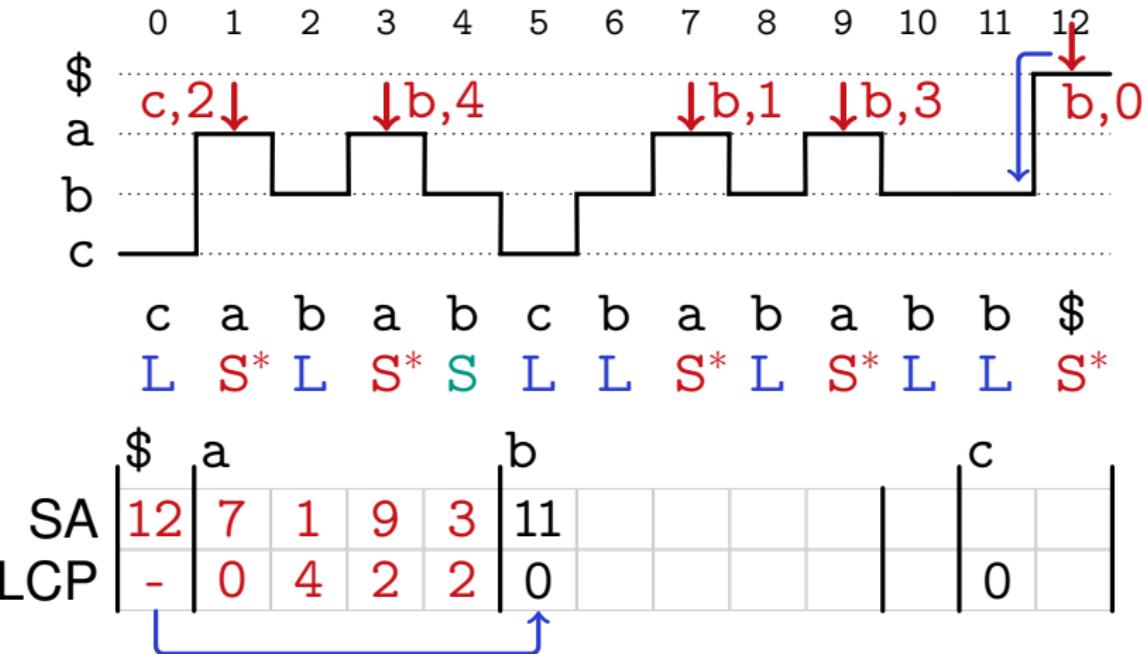
Example $T = [\text{cababcbababb\$}]$



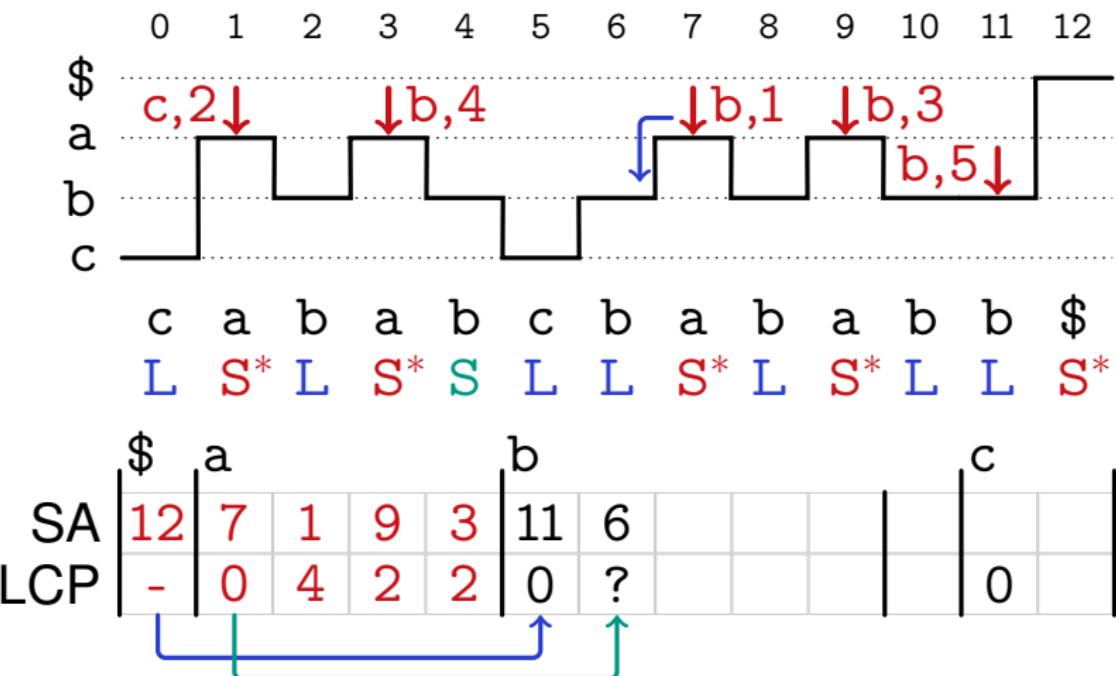
Example $T = [\text{cababcbababb\$}]$



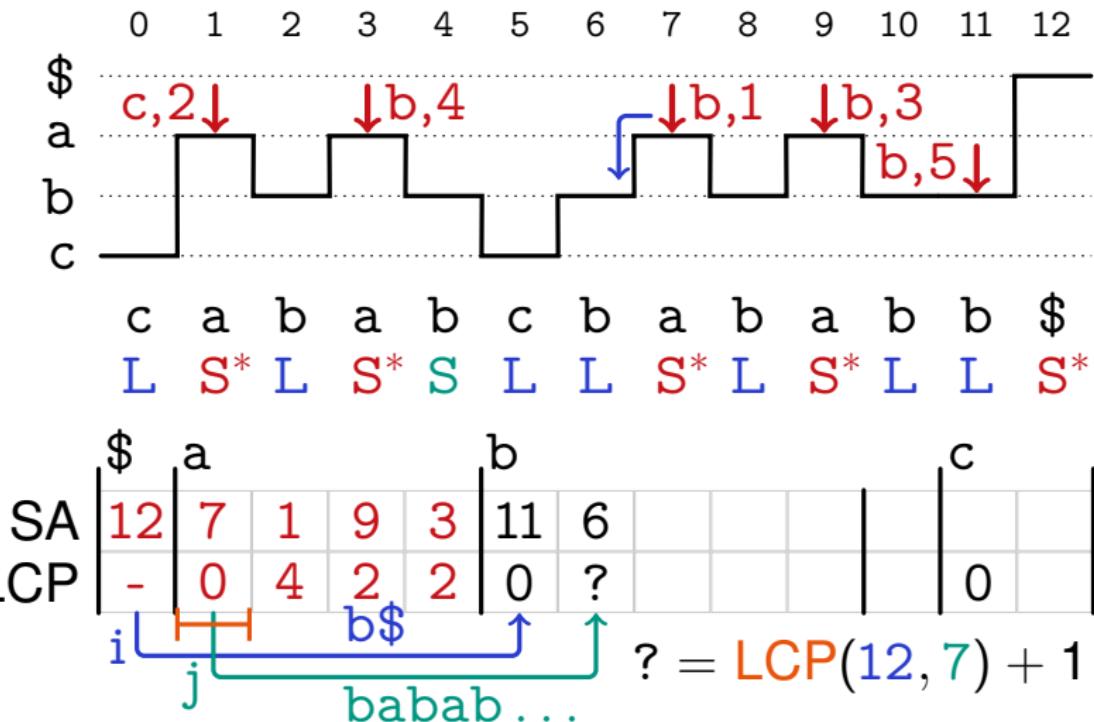
Example $T = [\text{cababcbababb\$}]$



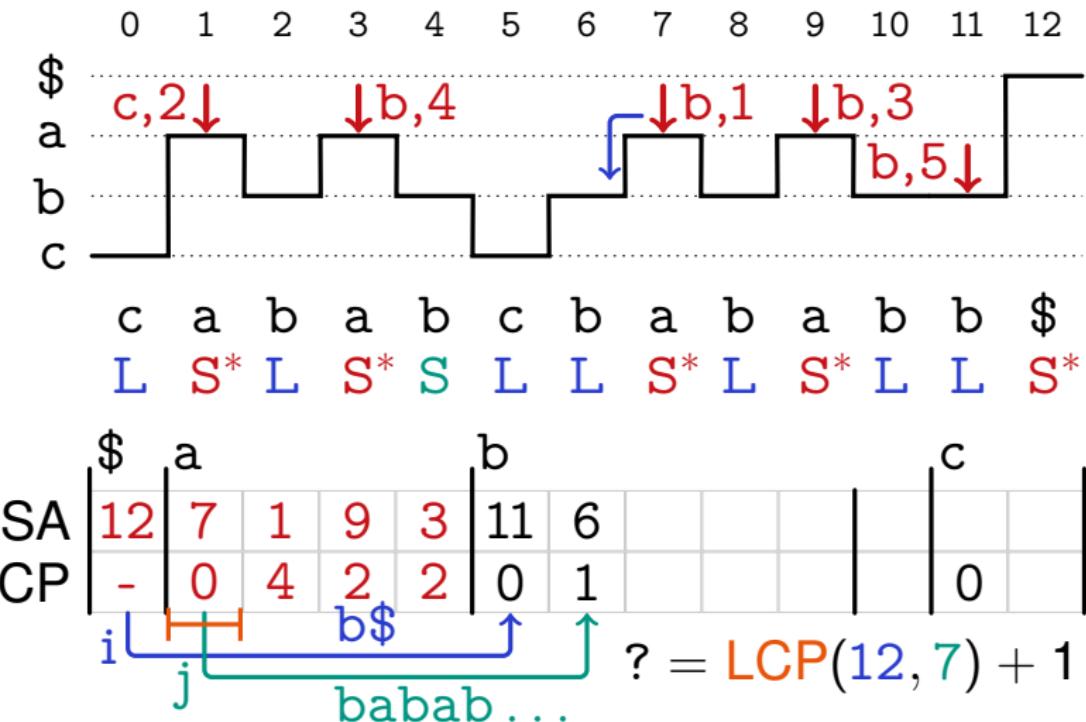
Example $T = [\text{cababcbababb\$}]$



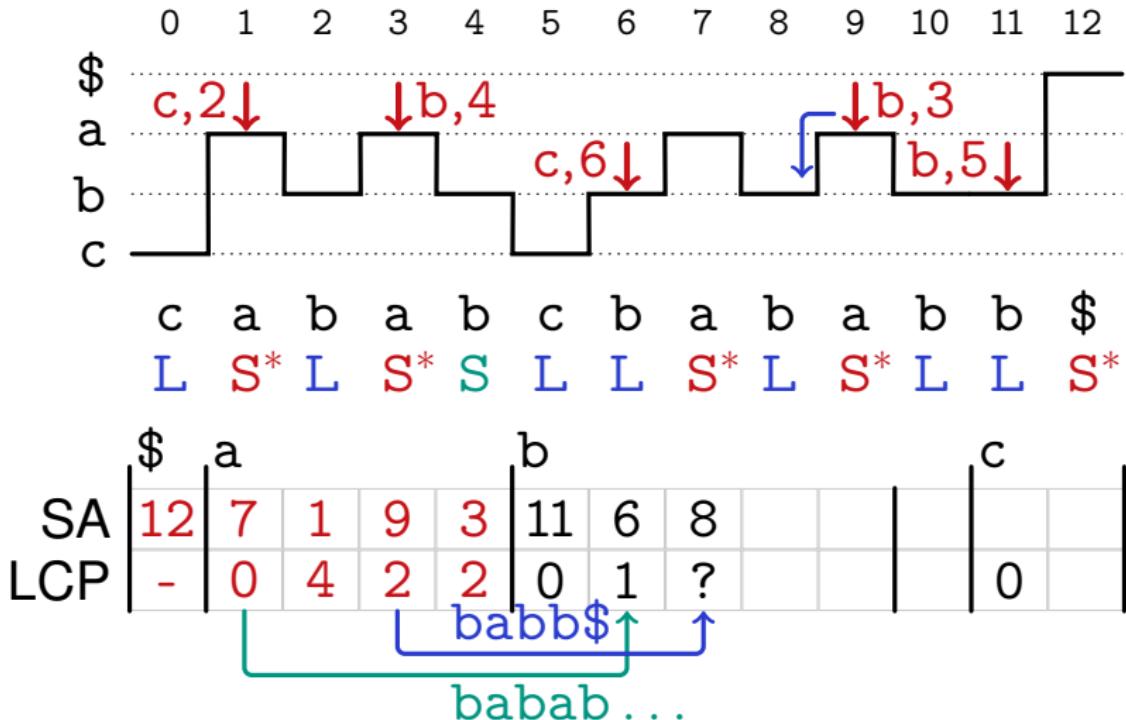
Example $T = [\text{cababcbababb\$}]$



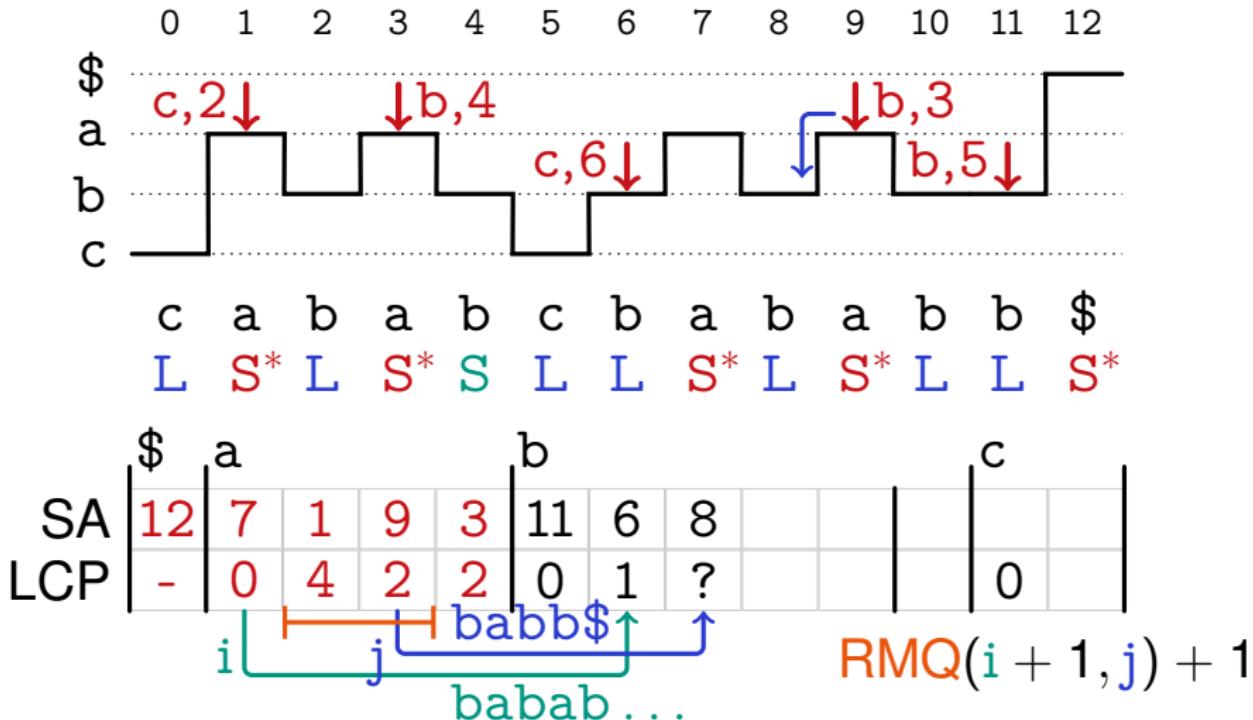
Example $T = [\text{cababcbababb\$}]$



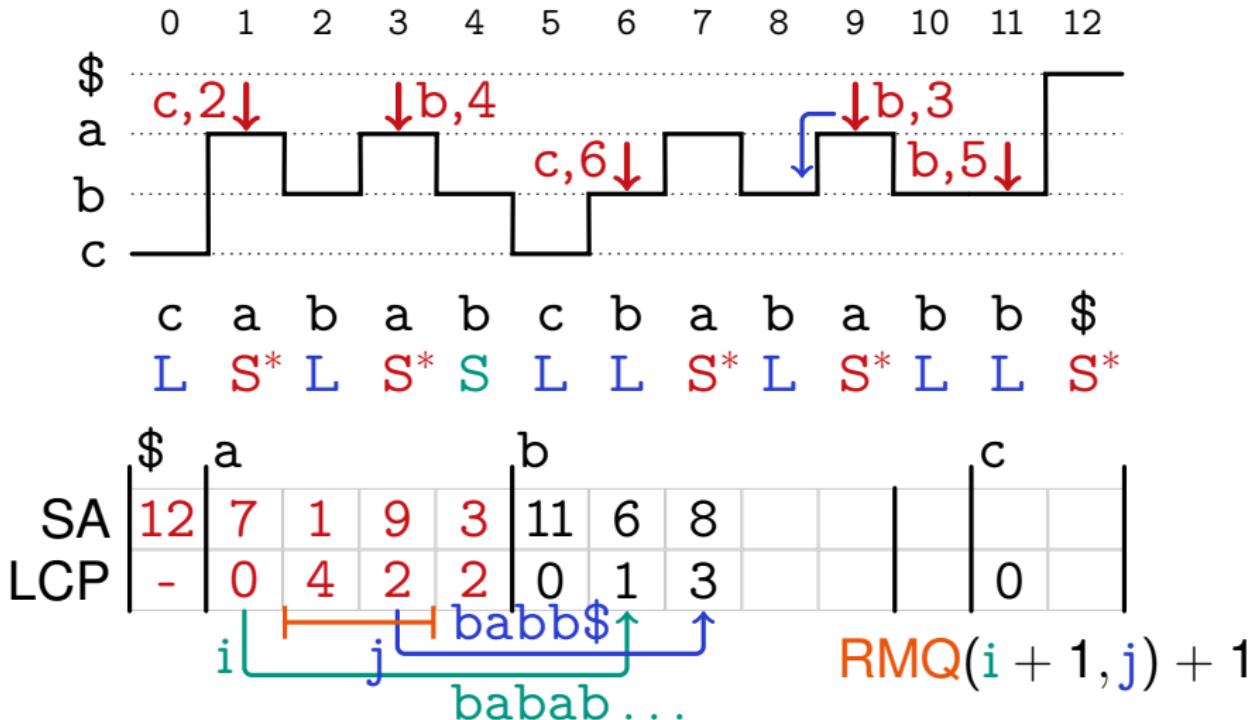
Example $T = [\text{cababcbababb\$}]$



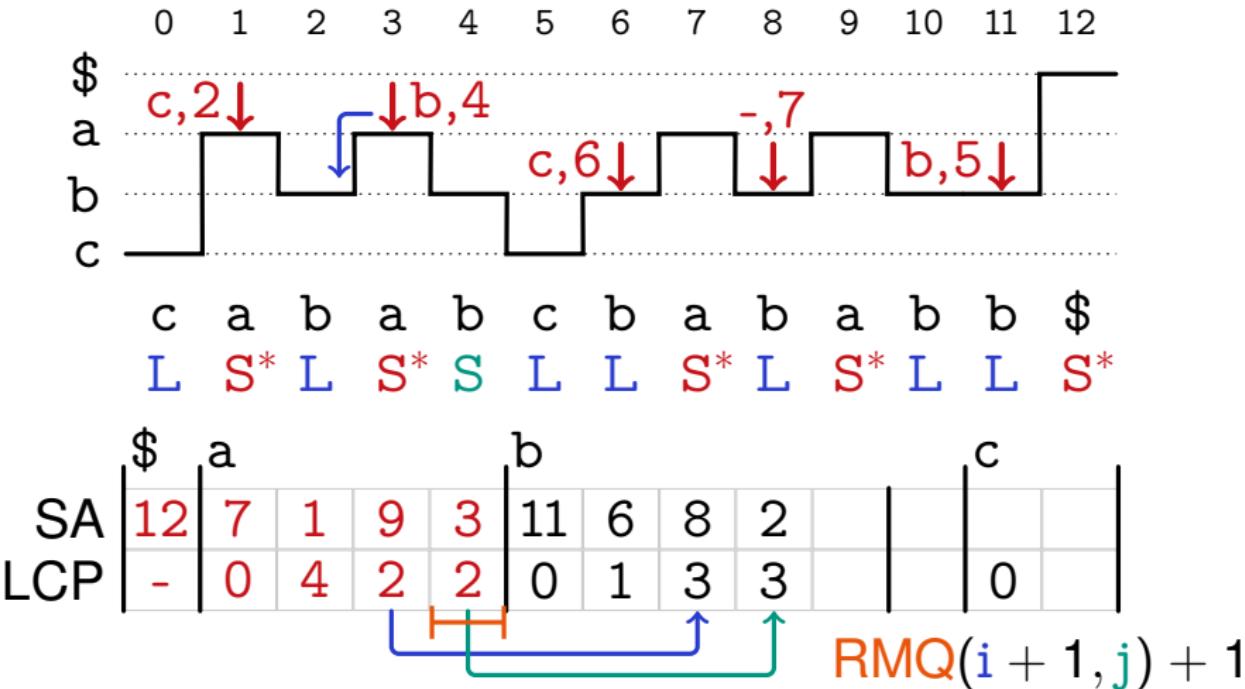
Example $T = [\text{cababcbababb\$}]$



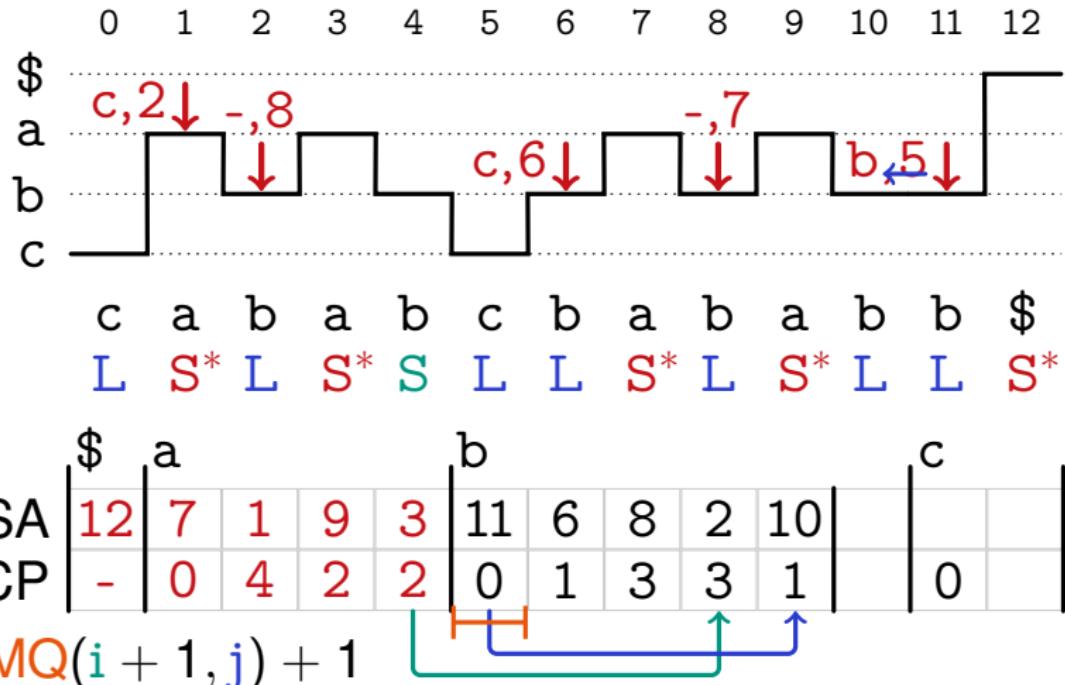
Example $T = [\text{cababcbababb\$}]$



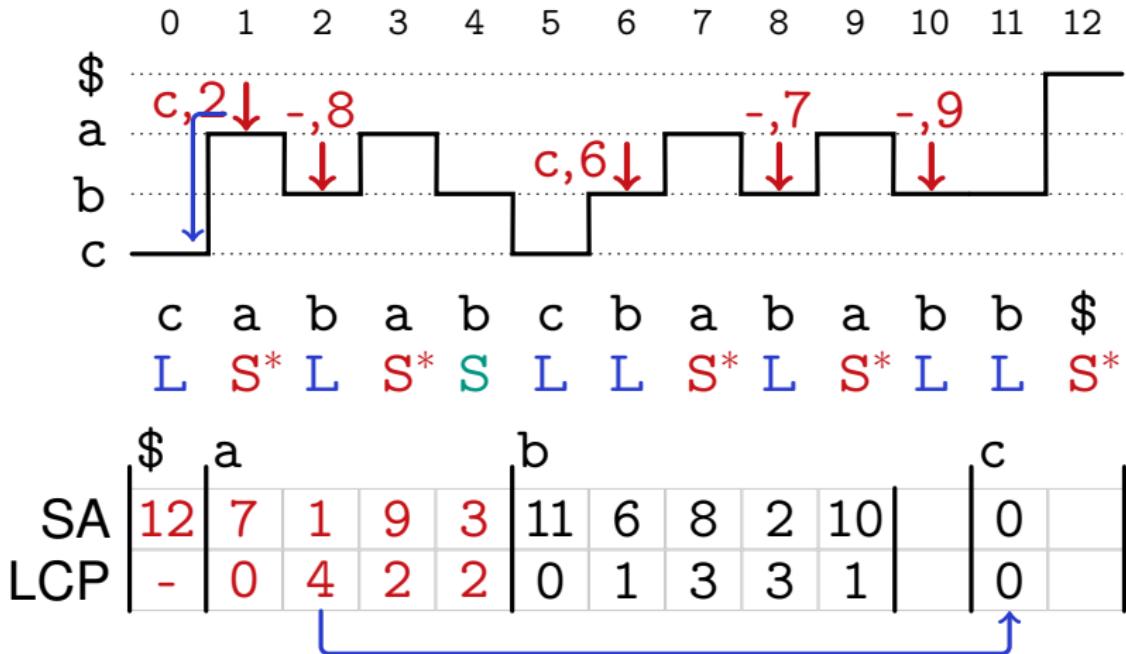
Example $T = [\text{cababcbababb\$}]$



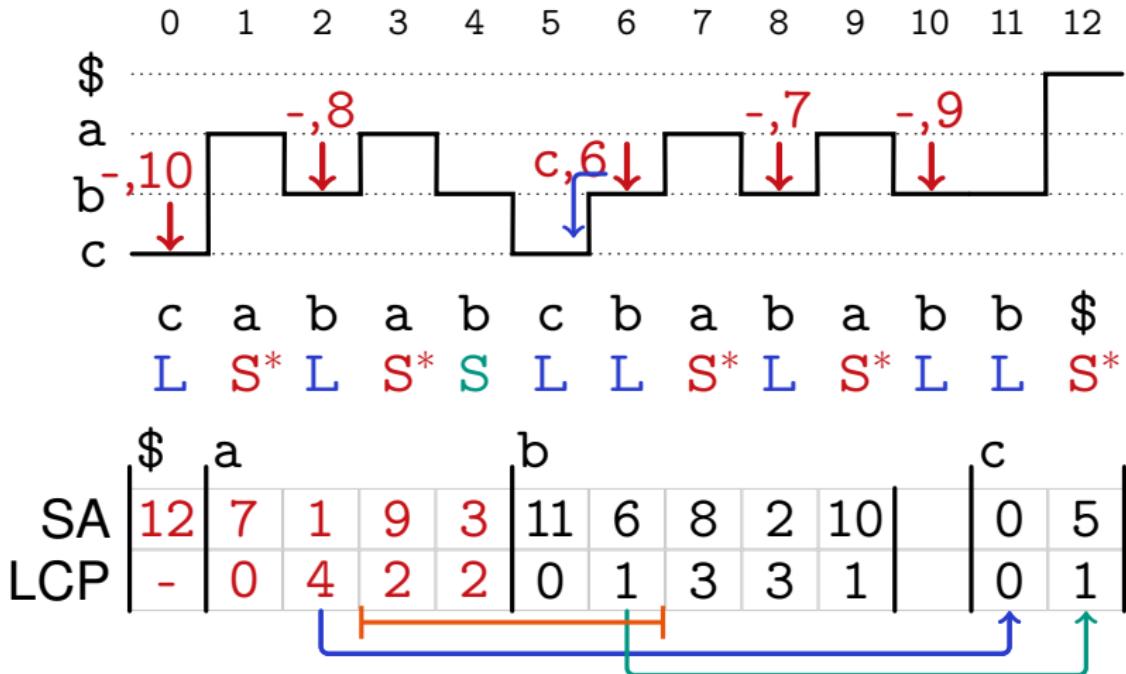
Example $T = [\text{cababcbababb\$}]$



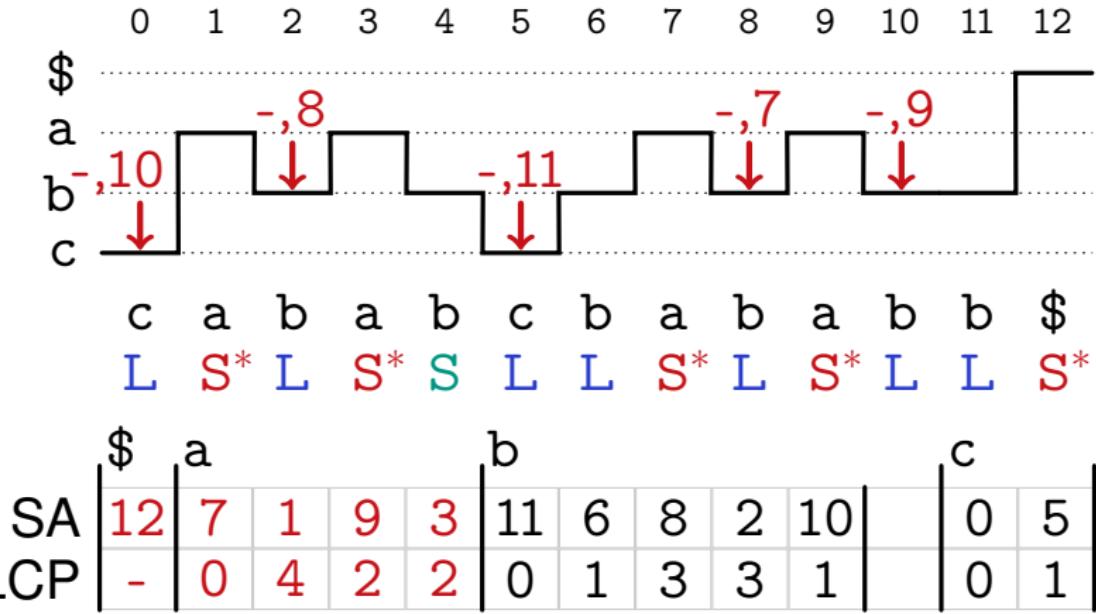
Example $T = [\text{cababcbababb\$}]$



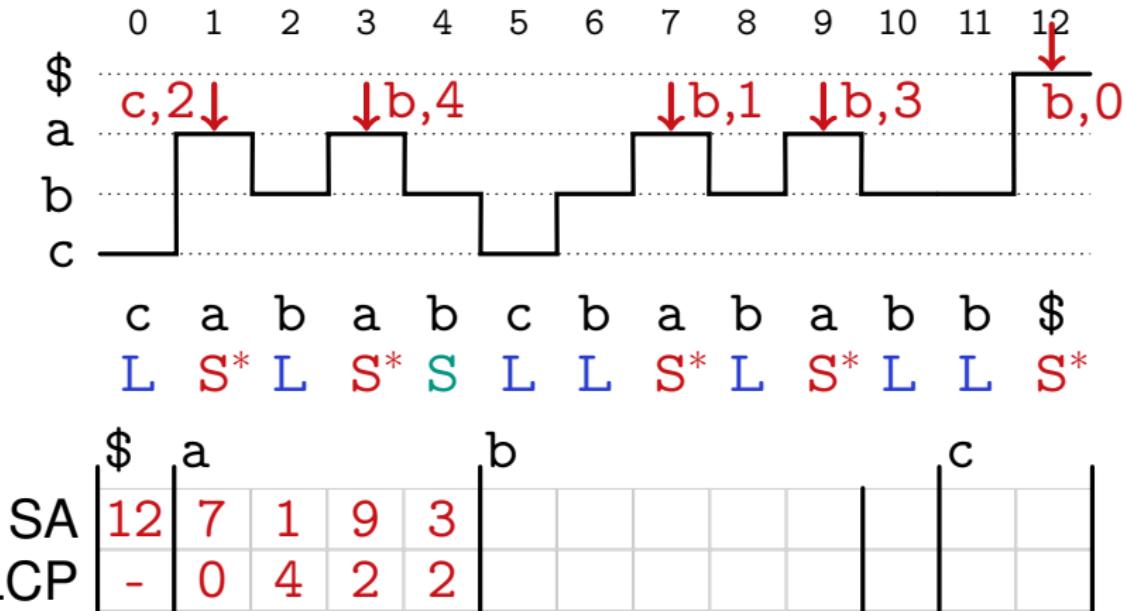
Example $T = [\text{cababcbababb\$}]$



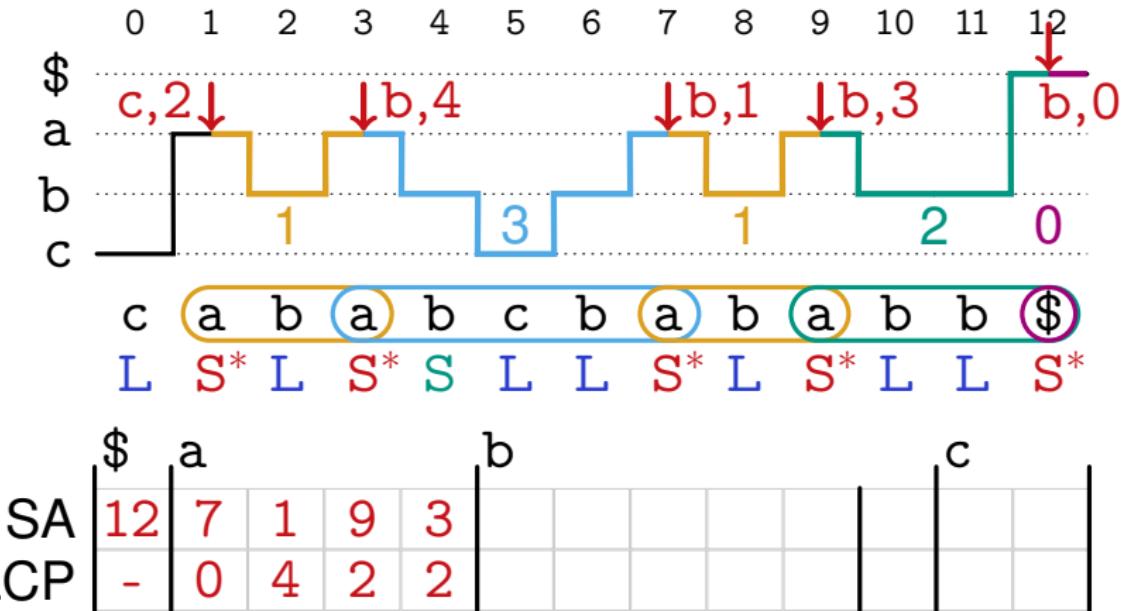
Example $T = [\text{cababcbababb\$}]$



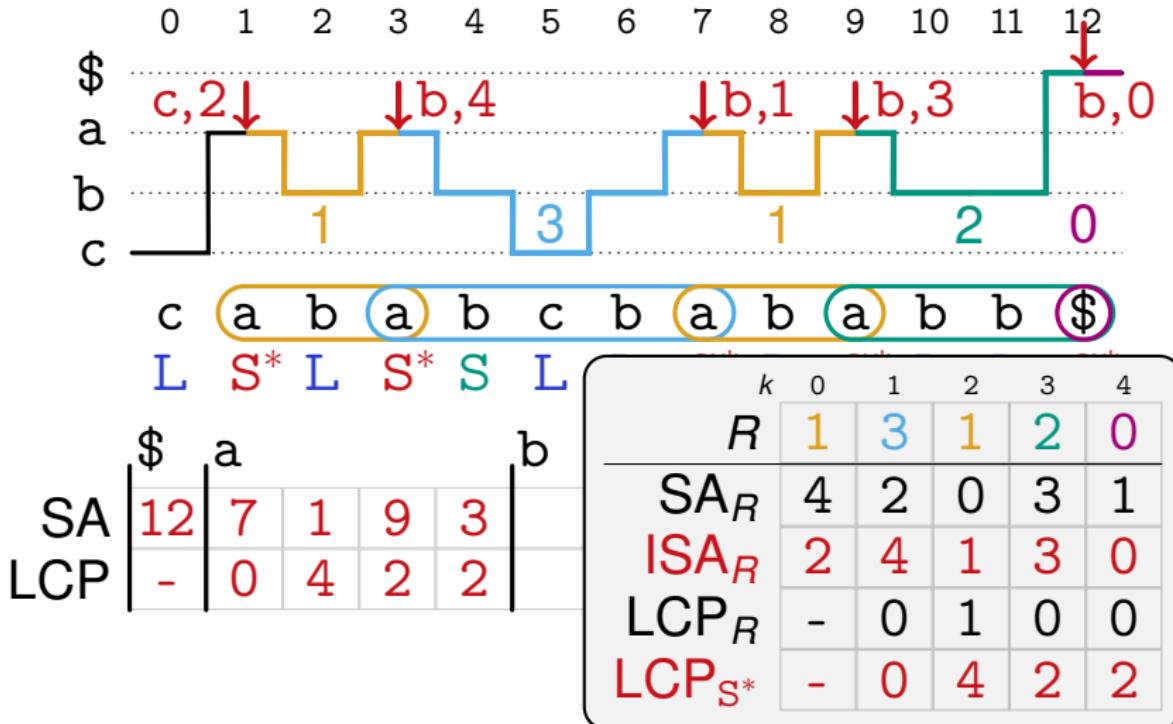
Example $T = [\text{cababcbababb\$}]$



Example $T = [\text{cababcbababb\$}]$



Example $T = [\text{cababcbababb\$}]$

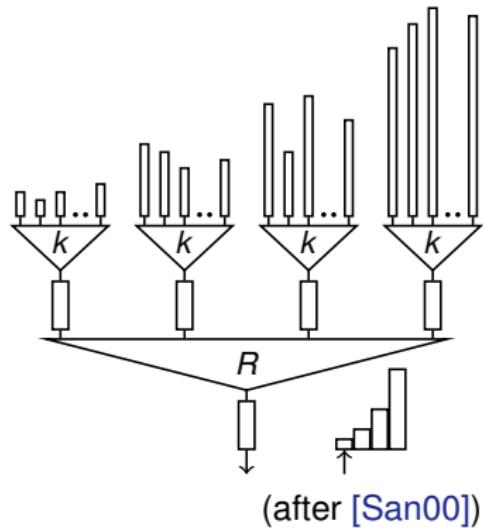


eSAIS with LCP Construction

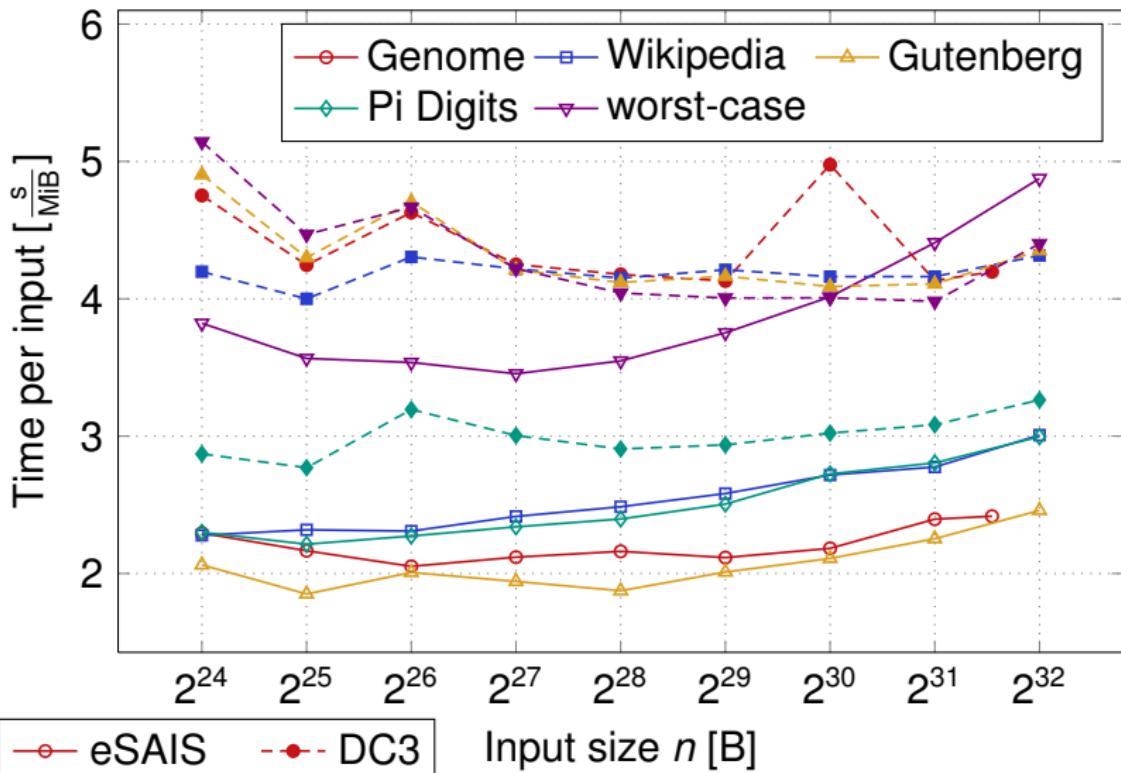
- Recursion size at most $\frac{n}{2}$.
- All SA construction steps: EM scanning, sorting, and PQ operations \Rightarrow amortized sorting complexity.
- LCP construction involves semi-dynamic RMQs.
Two solutions:
 - EM algorithm in sorting complexity for $N \leq M^2$, or
 - succinct in-memory data structure for practice.

Implementation Highlights

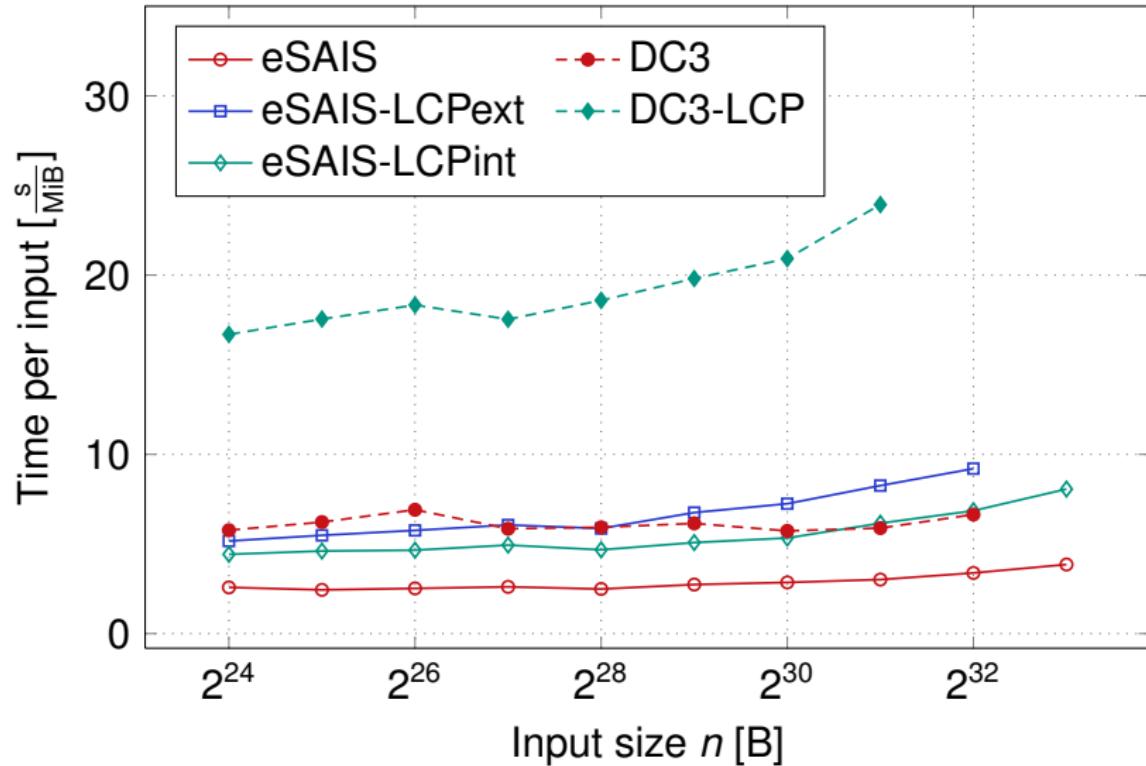
- Implemented in C++ using STXXL.
- STXXL [DKS08] provides efficient EM sorting and a priority queue.
- Enable 40-bit string positions.
- Further EM issues: handle large S^* -substrings via splitting.
- Source code available under GPL:
<http://tbingmann.de/2012/esais>



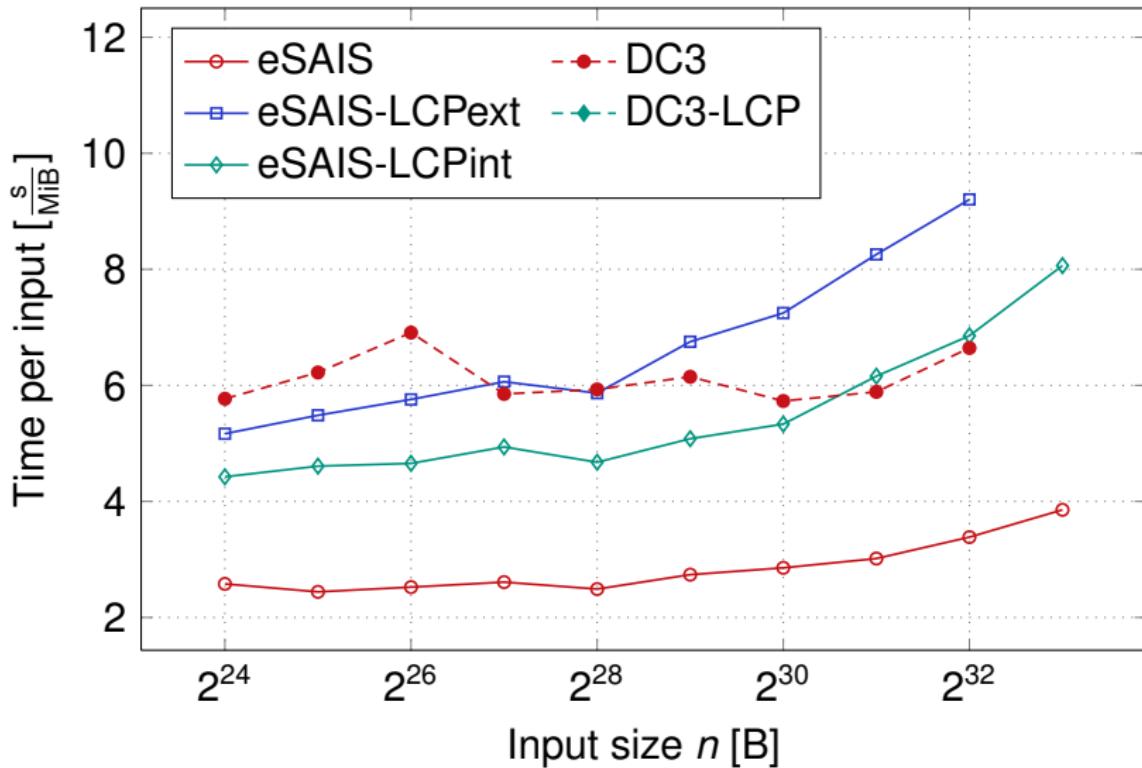
Construction SA only: eSAIS vs. DC3



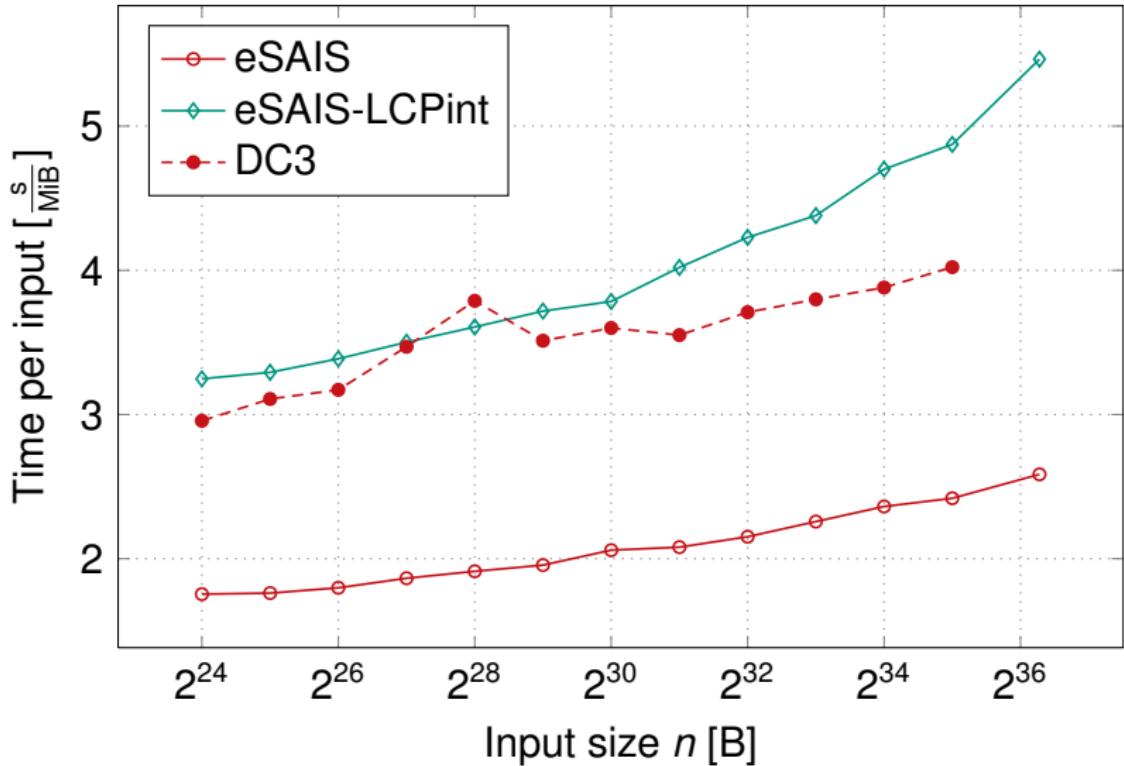
Construction Time: Gutenberg Text



Construction Time: Gutenberg Text



Construction Time: Wikipedia XML



Perspective?

Perspective?

Thank you for your attention!

Questions?

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